



OST

Eastern Switzerland
University of Applied Sciences

Blockchain (BlCh)

Repetition DSy – part 1

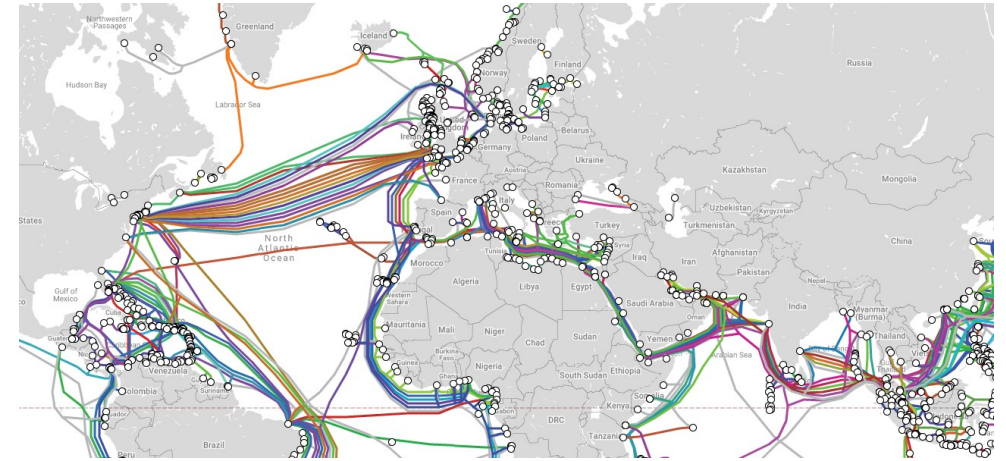
Thomas Bocek

17.09.2023

Lecture 1

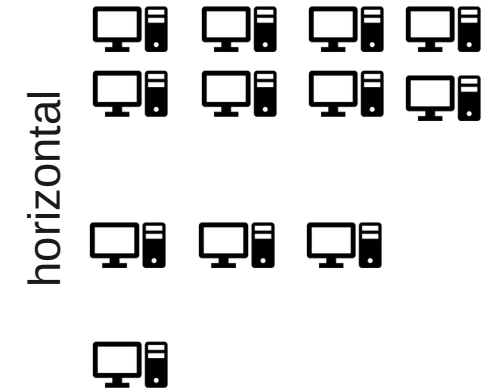
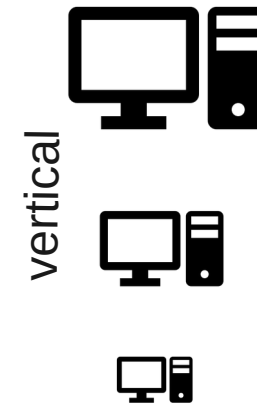
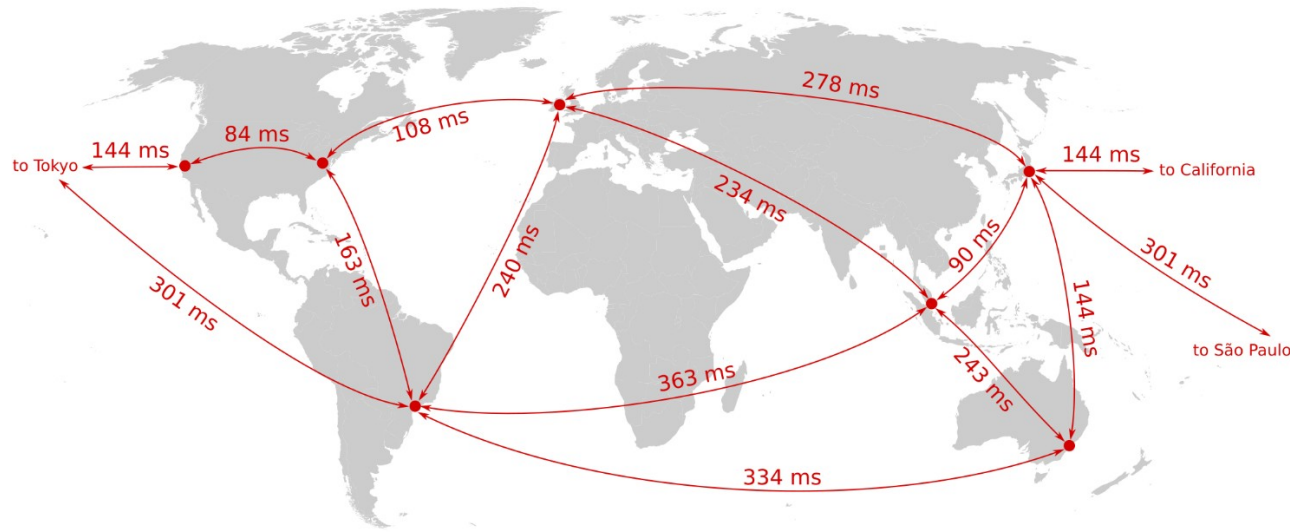
Distributed Systems Motivation

- Why Distributed Systems
 - **Scaling**
 - **Location**
 - **Fault-tolerance (bitflips, outages)**



Submarine Cable Map

<https://www.inkandswitch.com/local-first.html>



Lecture 2

Distributed Systems Categorization

“Controlled” Distributed Systems

- 1 responsible organization
- Low churn
- Examples:
 - Amazon DynamoDB
 - Client/server
- “Secure environment”
- High availability
- Can be homogeneous / heterogeneous

“Fully” Decentralized Systems

- N responsible organizations
- High churn
- Examples:
 - BitTorrent
 - Blockchain
- “Hostile environment”
- Unpredictable availability
- Is heterogeneous

Distributed Systems Categorization

“Controlled” Distributed Systems

- Mechanisms that work well:
 - Consistent hashing (DynamoDB, Cassandra)
 - Master nodes, central coordinator
- Network is under control or client/server → no NAT issues

“Fully” Decentralized Systems

- Mechanisms that work well:
 - Consistent hashing (DHTs)
 - Flooding/broadcasting - Bitcoin
- NAT and direct connectivity huge problem

Distributed Systems Categorization

“Controlled” Distributed Systems

- Consistency
 - Leader election (Zookeeper, Paxos, Raft)
- Replication principles
 - More replicas: higher availability, higher reliability, higher performance, better scalability, but: requires maintaining consistency in replicas
- Transparency principles apply

“Fully” Decentralized Systems

- Consistency
 - Weak consistency: DHTs
 - Nakamoto consensus (aka proof of work)
 - Proof of stake – Leader election, PBFT protocols
 - Is Bitcoin eventually consistent?
 - Some argue no, some argue it has even stronger guarantees [\[link\]](#)
- Replication principles apply to fully decentralized systems as well
- Transparency principles apply

Distributed Systems Categorization

- Spring Term – Distributed Systems (DSy)
 - Tightly/loosely coupled
 - Heterogeneous systems
 - Small-scale systems
 - Distributed systems
- Fall Term – Blockchain (BlCh)
 - Loosely coupled
 - Heterogeneous systems
 - Large-scale systems
 - Decentralized systems

(we will also talk about blockchains in this lecture)

(we will also talk about distributed systems in this lecture, but DSy is highly recommended)

Lecture 3

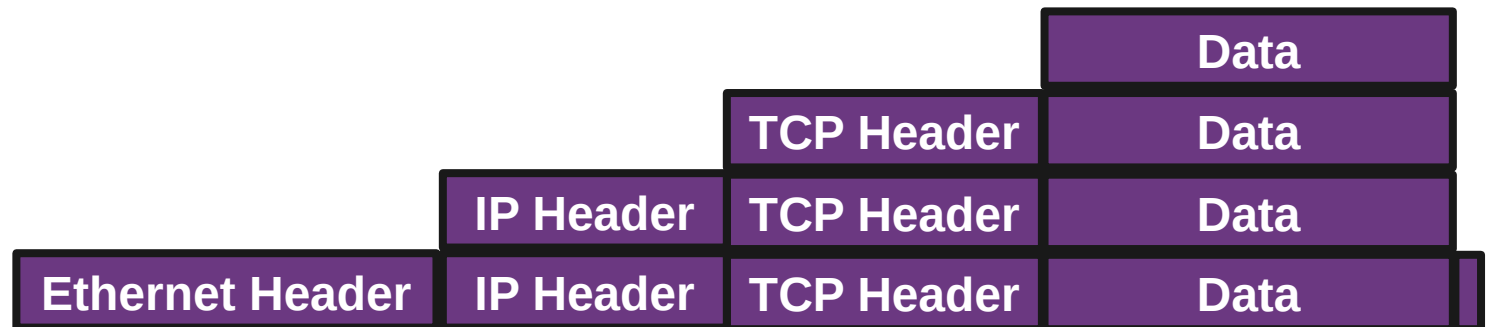
Pro/Cons - Opinion

- Monorepo
 - Tight coupling of projects
 - E.g., generating openapi.yml from backend, generate types for frontend → simply copy
 - Everyone sees all code / commits
 - Encourages code sharing within organization
 - Scaling: large repos, specialized tooling
- Opinion: [Accenture](#) - “From my experience, for a smaller team, starting with mono-repo is always safe and easy to start. Large and distributed teams would benefit more from poly-repo”
- My opinion: for small teams and “independent” project, use polyrepo. (I worked with small teams with mono and polyrepo, I have worked in big projects with polyrepos, but never in a big project with monorepos). If you have a tight coupling between projects (OpenAPI), use monorepos.
- Other opinion (sales pitch): <https://monorepo.tools>
- Polyrepo
 - Loose coupling of projects
 - If you want to generate openapi.yml, you need access from the backend repository to the frontend (e.g., curl+token)
 - Fine grained access control
 - Encourages code sharing across organizations
 - Scaling: many projects, special coordination

Networking: Layers

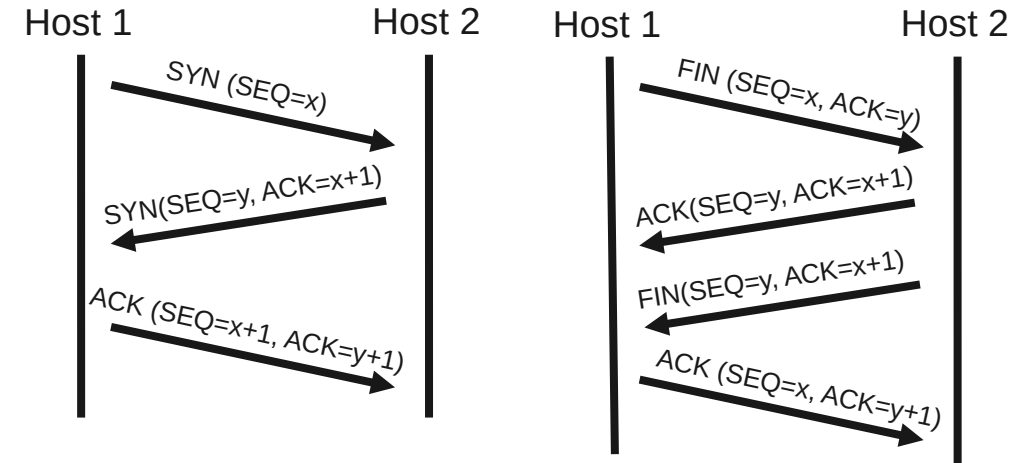
- Networking: Each vendor had its own proprietary solution - not compatible with another solution
 - [IPX/SPX](#) – 1983, [AppleTalk](#) 1985, [DECnet](#) 1975, [XNS](#) 1977
- Nowadays most vendors build compatible networks hardware/software from different vendors
 - Cisco, Dell, HP, Huawei, Juniper, Lenovo, Linksys, Netgear, MicroTik, Siemens, Ubiquiti, etc.
- Goal of layers: interoperability
 - 1984: ISO 7498 - The Basic Reference Model for Open Systems Interconnection

OSI model	"Internet model"
Application	Application
Presentation	
Session	
Transport	Transport
Network	Internet
Data link	Link
Physical	



Layer 4 - TCP

- Connection establishment
 - SYN, SYN-ACK, ACK (three way)
 - Initiates TCP session: initial sequence number is ~ random
- Connection termination
 - FIN, ACK + FIN, ACK (three/four way)
 - 3-way handshake, when host 1 sends a FIN and host 2 replies with a FIN & ACK
- Sequences and ACKs
 - Identification each byte of data
 - Order of the bytes → reconstruction
 - Detecting lost data: RTO, DupACK:



- Retransmission timeout
 - If no ACK is received after timeout (e.g. 2xRTT), resend.
- Duplicate cumulative acknowledgements, selective ACK [[link](#)]
 - ACKs for last consecutive packets
 - 3 times same ACK → retransmit missing packets (fast retransmit)

TCP/IP from an Application Developer View

- Server in golang ([repo](#))
 - git clone
<https://github.com/tbocek/DSy>
 - Download [GoLand](#), or [others](#)
 - go run server.go → server
- Listening on TCP port 8081
 - Return string in uppercase
- Node.js version
 - Download [WebStorm](#), or [other](#)
- Client:
 - nc localhost 8081

```
const net = require('net');
const server = new net.Server();
server.listen(8081, function() {
  console.log('Launching server...');
});

server.on('connection', function(socket) {
  socket.on('data', function(chunk) {
    console.log('Data received from client: $
{chunk.toString()}');

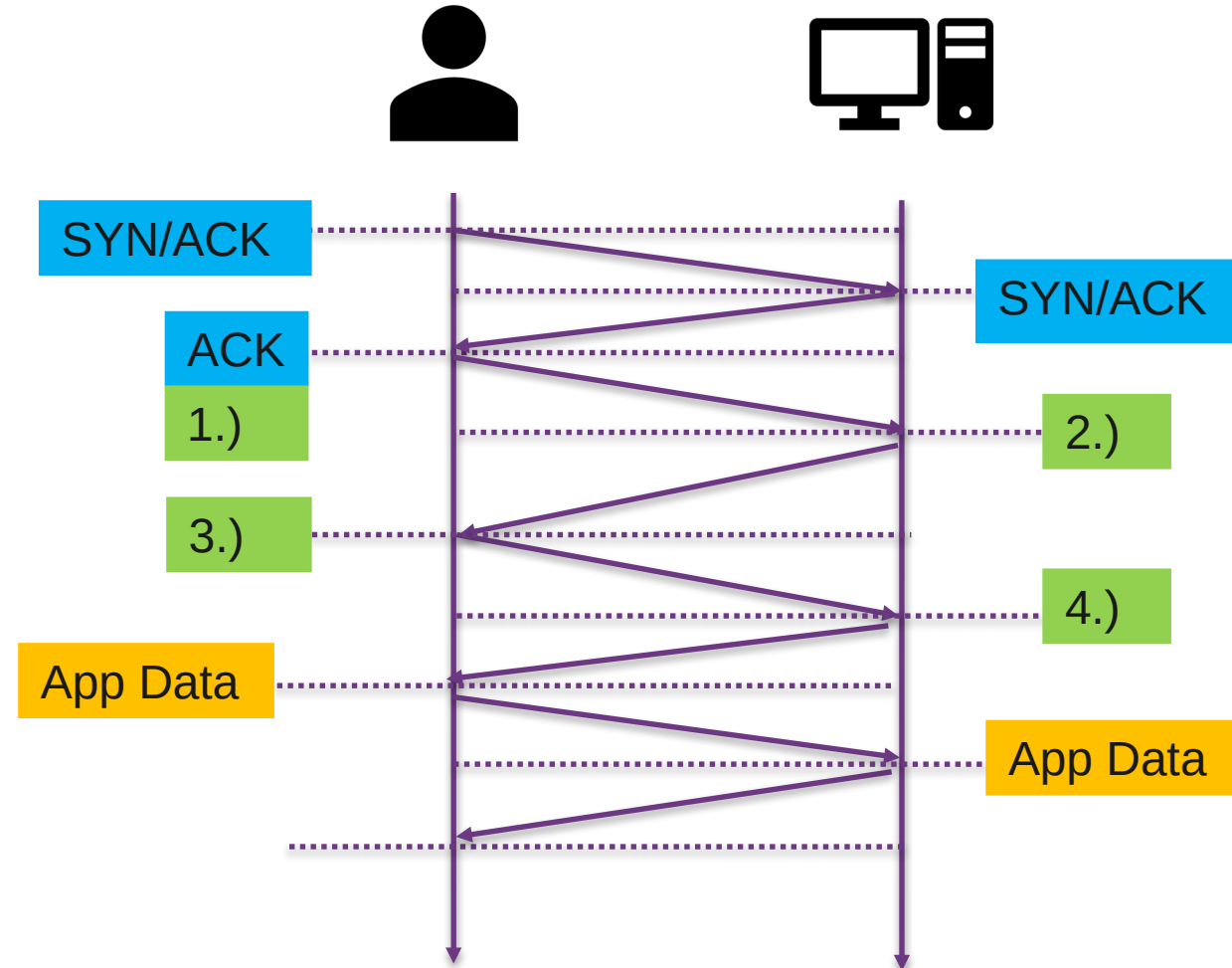
socket.write(chunk.toString().toUpperCase() +
"\n");
  });
});
```

```
package main
import ("bufio"
  "fmt"
  "net"
  "strings")
func main() {
  fmt.Println("Launching server...")
  ln, _ := net.Listen("tcp", ":8081") // listen
on all interfaces
  for {
    conn, _ := ln.Accept() // accept
connection on port
    message, _ :=
bufio.NewReader(conn).ReadString('\n') //read line
    fmt.Print("Message Received:",
string(message))
    newMessage := strings.ToUpper(message)
//change to upper
    conn.Write([]byte(newMessage + "\n"))
//send upper string back
  }
}
```

Lecture 4

Layer 4 – TCP + TLS

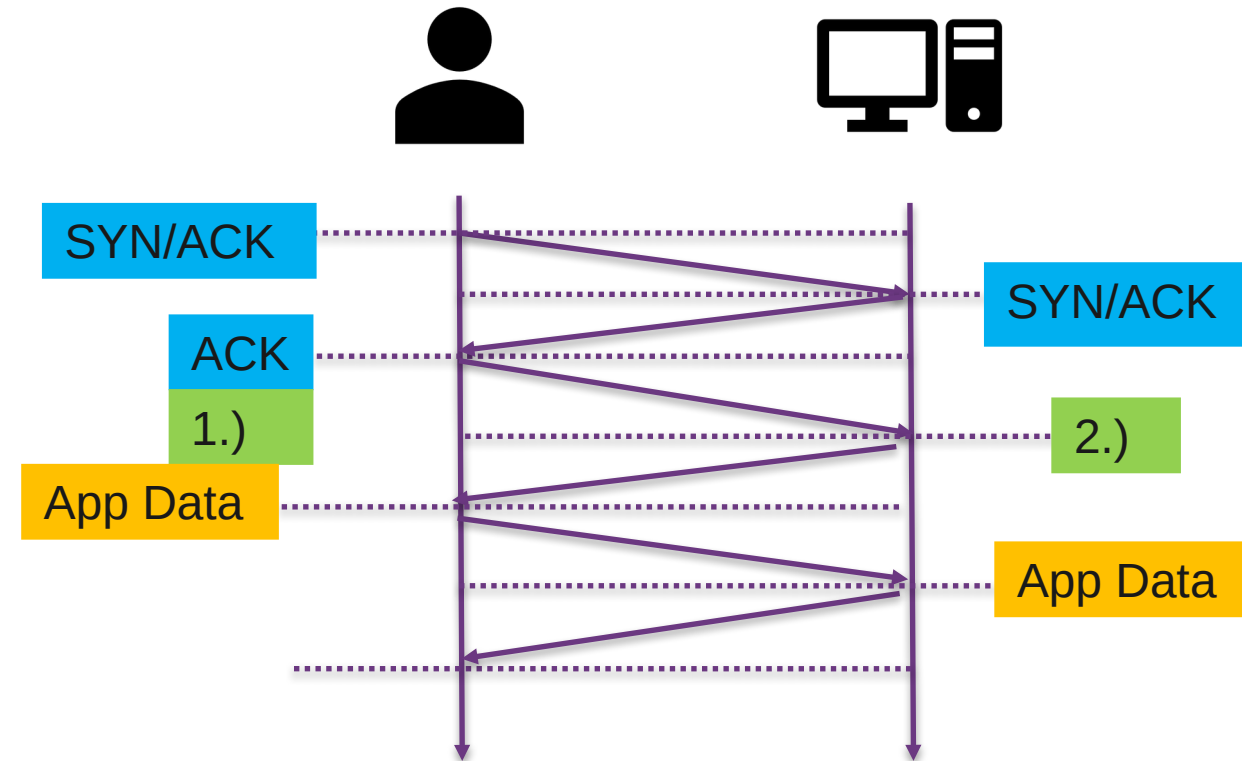
- Security: Transport Layer Security (TLS)
 1. "client hello" lists cryptographic information, TLS version, ciphers/keys
 2. "server hello" chosen cipher, the session ID, random bytes, digital certificate (checked by client), optional: "client certificate request"
 3. Key exchange using random bytes, now server and client can calc secret key
 4. "finished" message, encrypted with the secret key
- 3 RTT to send first byte, 4RTT to receive first byte




```
PING sydney.edu.au (129.78.5.8) 56(84) bytes of data.  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=1 ttl=233 time=307 ms  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=2 ttl=233 time=305 ms  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=3 ttl=233 time=305 ms
```

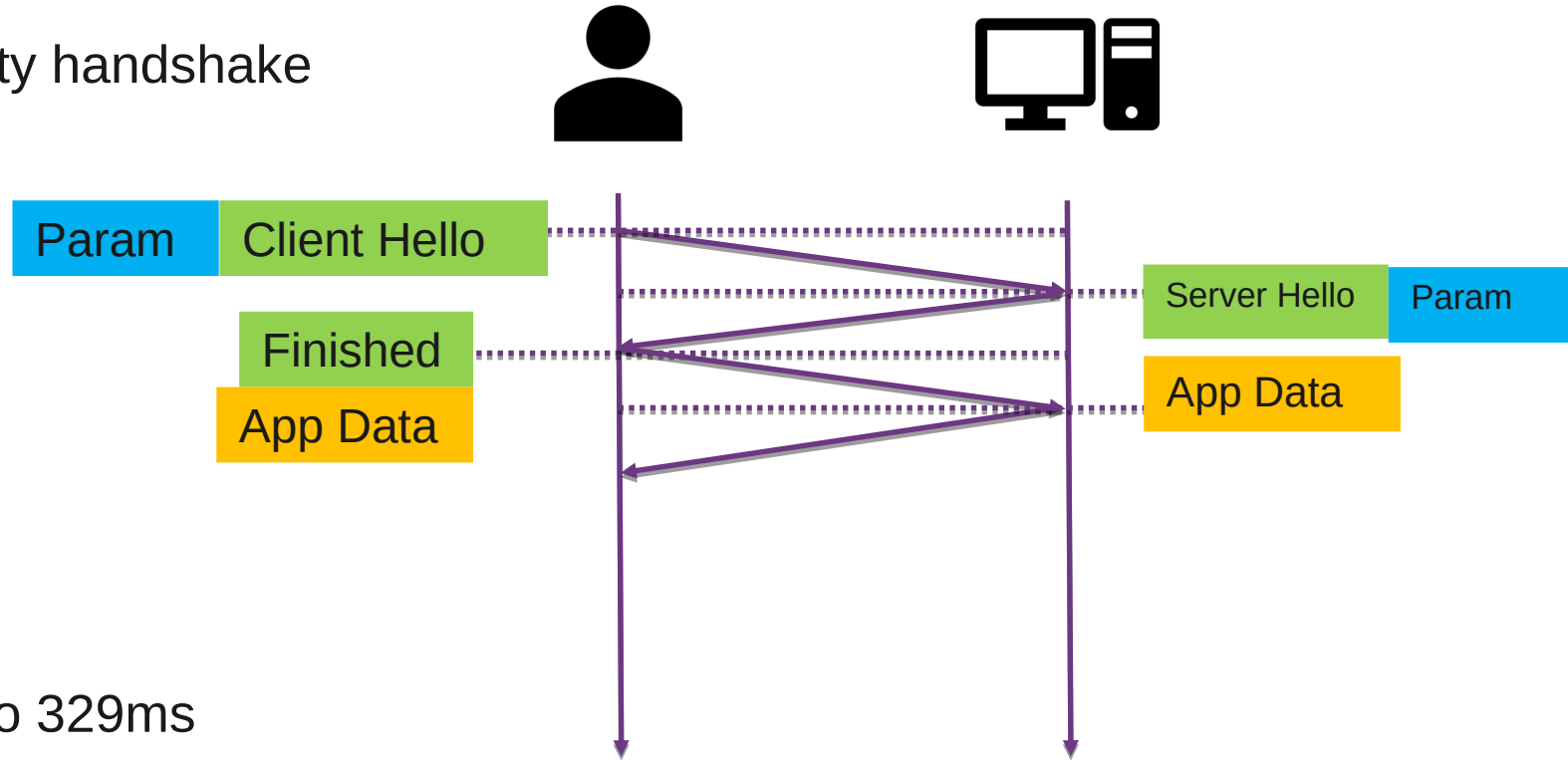
Layer 4 – TCP + TLS

- Ping to Australia: 329ms
 - One way ~ 165ms
- TCP + TLS handshake:
 - 3RTT = 987ms! No data sent yet
- TLS 1.3, finished Aug 2018
 - 1 RTT instead of 2
 - 1.) Client Hello, Key Share
 - 2.) Server Hello, key Share, Verify Certificate, Finished
 - 0 RTT possible, for previous connections, losing perfect forward secrecy
- 90% of browsers used already support it



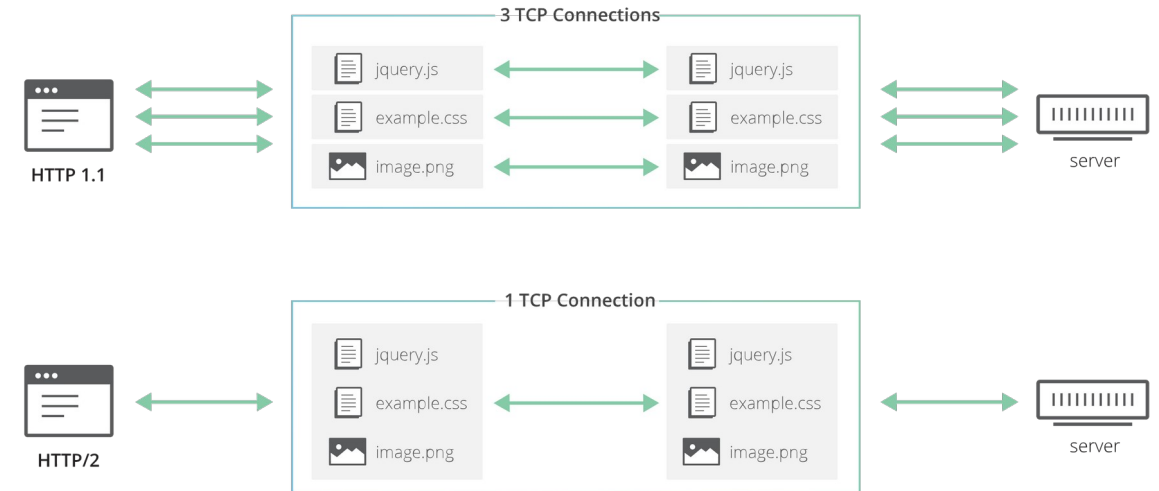
QUIC / HTTP/3

- QUIC: 1RTT connection + security handshake
 - For known connections: 0RTT
 - Built in security
 - “Google's 'QUIC' TCP alternative slow to excite anyone outside Google” [link] ([9%](#), [25%](#), 75%)
 - [Facebook](#)
 - [Cloudflare](#), [state of HTTP](#)
- Example Australia: from 987ms to 329ms



QUIC / HTTP3

- Multiplexing in HTTP/2
 - [HTTP/1 → HTTP/2](#)
- HTTP/2: Head-of-line blocking
 - One packet loss, TCP needs to be ordered
 - QUIC can multiplex requests: one stream does not affect others
- HTTP/3 is great, but...
 - NAT → SYN, ACK, FIN, conntrack knows when connection ends, not with QUIC, timeouts, new entries, many entries
 - HTTP header compression, referencing previous headers
 - Many TCP [optimizations](#)



source: <https://blog.cloudflare.com/the-road-to-quic/>

