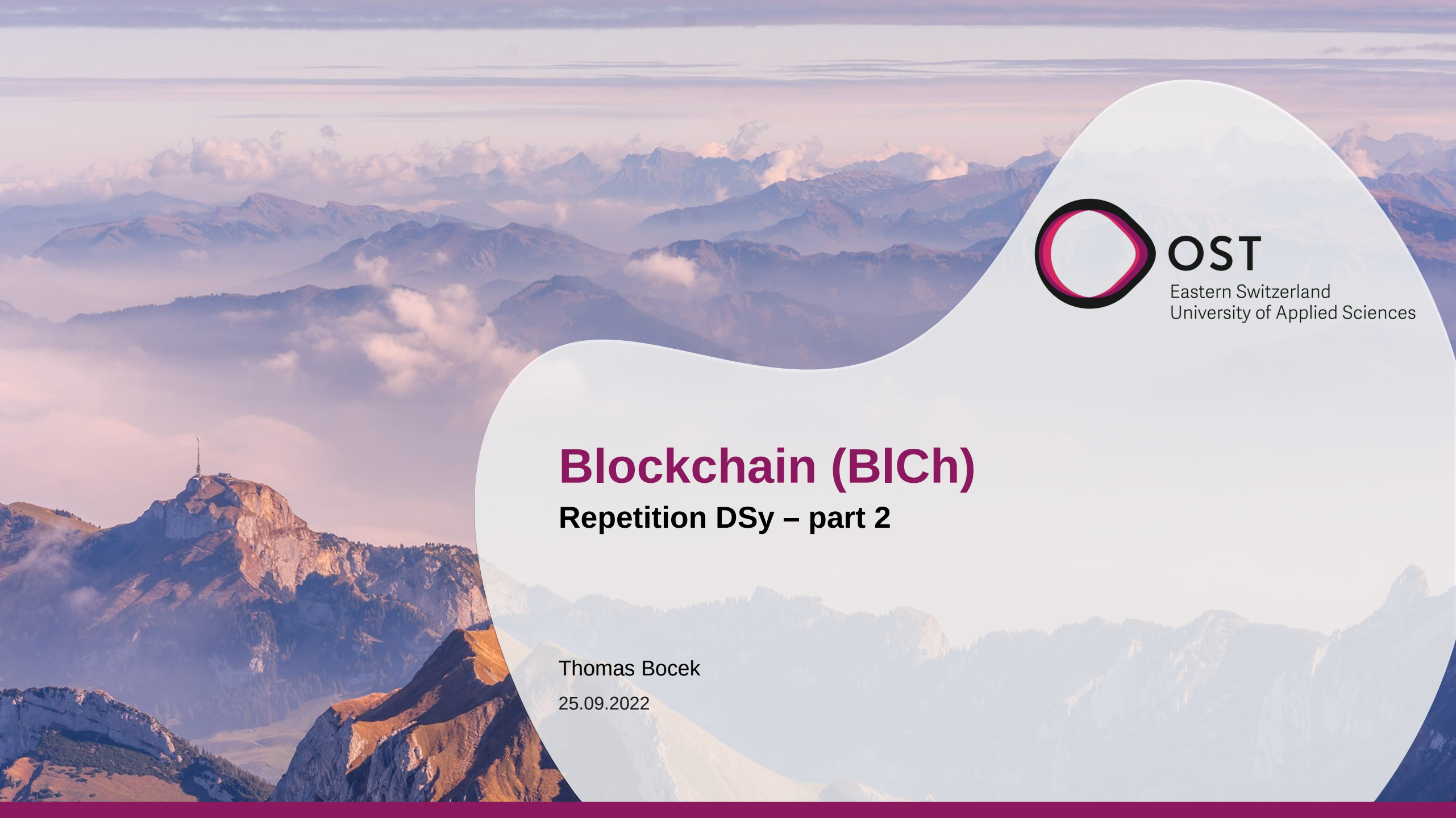

OST
Eastern Switzerland
University of Applied Sciences

Blockchain (BlCh)

Repetition DSy – part 2

Thomas Bocek

25.09.2022

Lecture 6

Authentication

- Authentication
  - Single-factor authentication
    - E.g. password
  - Multi-factor authentication / 2FA
    - E.g. password **and** software token, SMS (15.03.2021)
- Password rules
  - Don’t use:
    - The name of a pet, child, family member, or significant other
    - Anniversary dates and birthdays
    - Birthplace
    - Name of a favorite holiday
    - Something related to a favorite sports team
    - The word "password“
  - Don’t’ reuse passwords, use password managers

- Don’t enter passwords on unencrypted sites
- Password length:
  password cracking with 5000$ in 2018 with hashcat
  - Hashtype: WPA/WPA2: 1190.5 kH/s

Pw length	Combinations	Time
6	11m	9s
7	656m	9m
8	38b	8h
9	$7 * 10^{15}$	186y
10	$4 * 10^{17}$	11ky
11	$2 * 10^{19}$	665ky
12	$1 * 10^{21}$	38my

- Combinations depend on PW complexity

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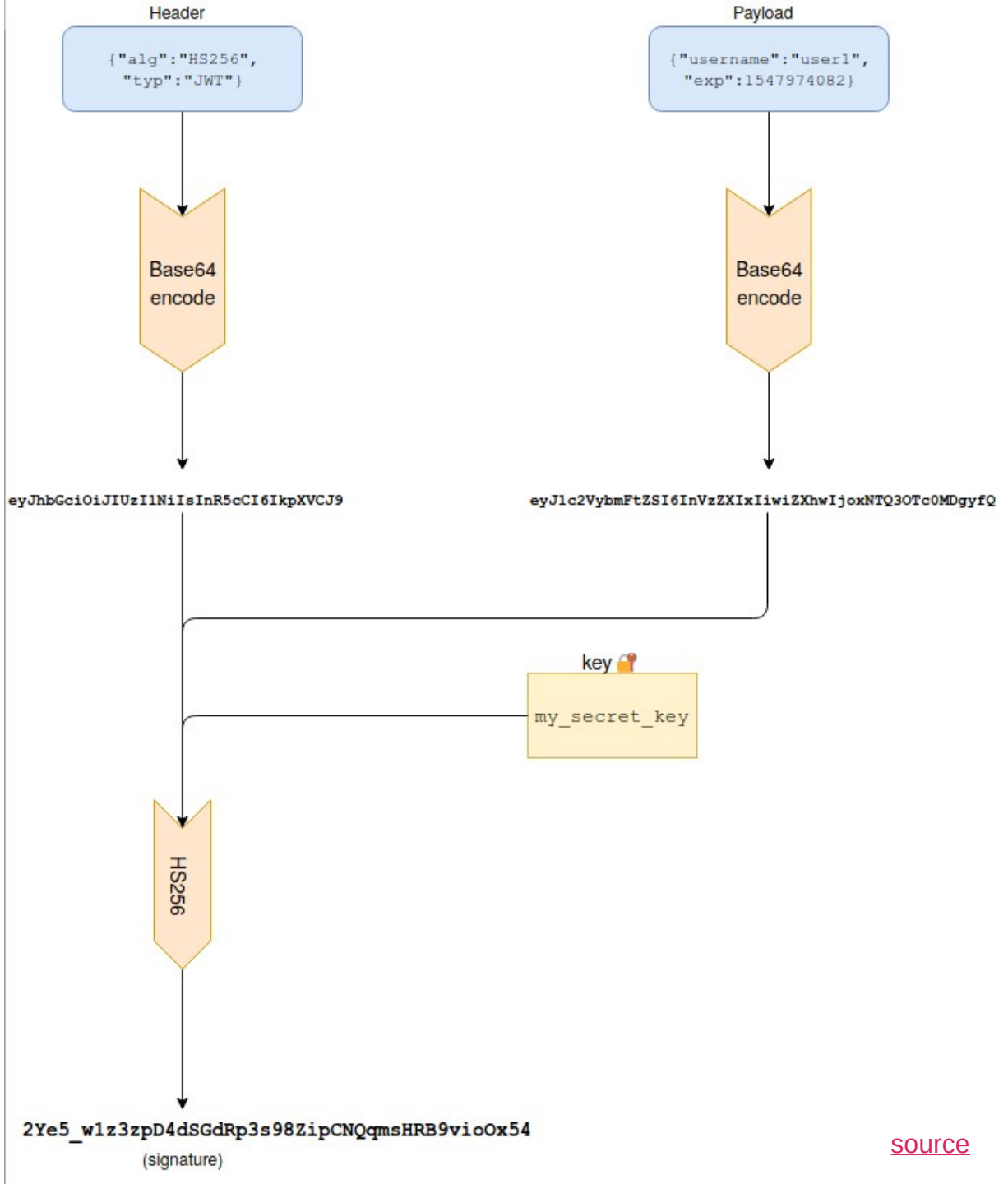
Authentication

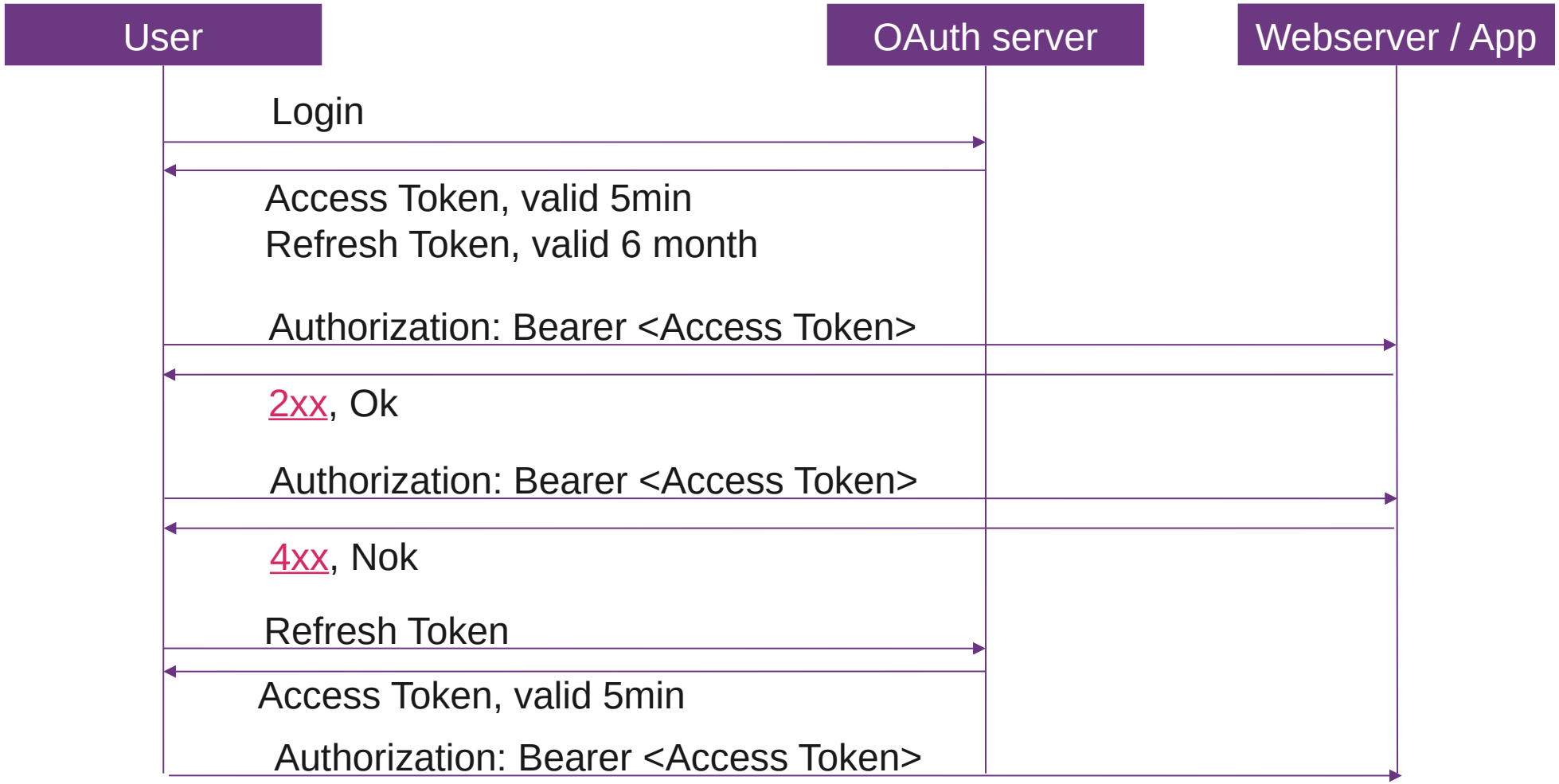
- JSON-based access tokens
  - Header: {"alg" : "HS256"}
  - Payload: {“sub" : "tom", "role" : "admin", “exp" : 1422779638}
- Signature (simple): keyed-hash message
  - ~hash(base64(header)+base64(payload) + secret token)
- Client can store user_token in
  - localStorage.setItem(“token", userToken);
- Example in golang with JWT
  - Tutorial: here and here

- OAuth - protocol for authorization 3rd party integration
  - Grant access on other websites without giving them the passwords



source

Access Token / Refresh Token



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Access Token / Refresh Token

- Access Token only short lifetime, e.g., 10min.
  - If public key / secret is known, the content in the token can be trusted, e.g., in the serivce
  - Can have userId, role, etc.
    - No need to query DB for those information, e.g.:

```
type TokenClaims struct {
      MailFrom string `json:"mail_from,omitempty"`
      MailTo   string `json:"mail_to,omitempty"`
      jwt.Claims
}
```

- Refresh Token longer lifetime, e.g., 6 month
  - A refresh token is used to get a new access token
  - IAM / Auth server creates access tokens

- Only access token, with long lifetime
  - If a user credential is revoked – how to inform every service?
- Only refresh token
  - Tightly coupled Service/Auth, every request to Service, Auth needs to be involved for every access
- Access + Refresh token
  - If a user credential is revoked, user has max. 10min more to access service
  - Auth only involved if access token is expired
- Authorization Code Flow with Proof Key for Code Exchange (PKCE)

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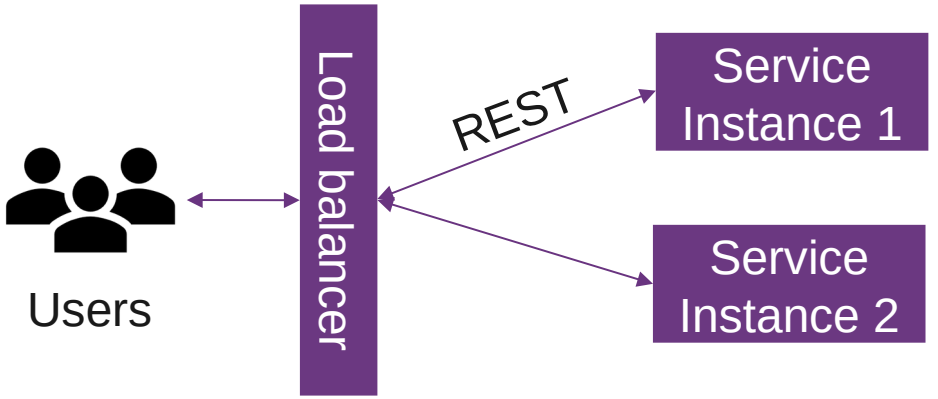
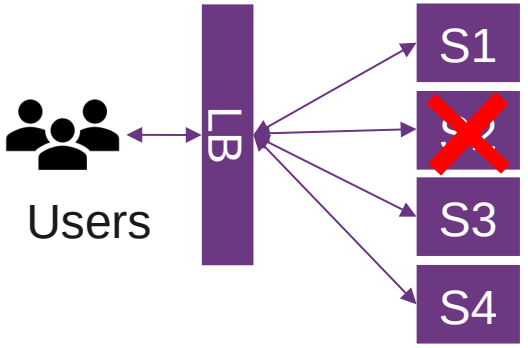
Load Balancing

- What is load balancing
  - Distribution of workloads across multiple computing resources
    - Workloads (requests)
    - Computing resources (machines)
  - Distributes client requests or network load efficiently across multiple servers [link]
    - E.g., service get popular, high load on service

→ horizontal scaling

- Why load balancing
  - Ensures high availability and reliability by sending requests only to servers that are online
  - Provides the flexibility to add or subtract servers as demand dictates





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Caddy

- Configuration: dynamic
  - Static: Caddyfile
- One-liners:
  - Quick, local file server: caddy file-server
  - Reverse proxy: caddy reverse-proxy --from example.com --to localhost:9000

```
:7070
reverse_proxy 127.0.0.1:8081 127.0.0.1:8080 {
  unhealthy_status 5xx
  fail_duration 5s
}
```

- Open Source, software-based load balancer: https://github.com/caddyserver/caddy
  - “Caddy 2 is a powerful, enterprise-ready, open source web server with automatic HTTPS written in Go”
  - L7 load balancer
  - Reverse proxy
  - Static file server
  - HTTP/1.1, HTTP/2, and experimental HTTP/3
  - Caddy on docker hub

Dockerfile

- Example: caddy as LB, go as Service
  - docker-compose up --scale services=5

```
#docker-compose.yml
version: '3'
services:
  services:
    build: .
    ports:
      - "8080-8085:8080"
  lb:
    image: caddy
    ports:
      - "7070:7070"
    volumes:
      - ./Caddyfile:/etc/caddy/Caddyfile
```

```
#Caddyfile
:7070
reverse_proxy * {
  to http://dsy-services-1:8080
  to http://dsy-services-2:8080
  to http://dsy-services-3:8080
  to http://dsy-services-4:8080
  to http://dsy-services-5:8080

  lb_policy round_robin
  lb_try_duration 1s
  lb_try_interval 100ms
  fail_duration 10s
  unhealthy_latency 1s
}
```


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CORS

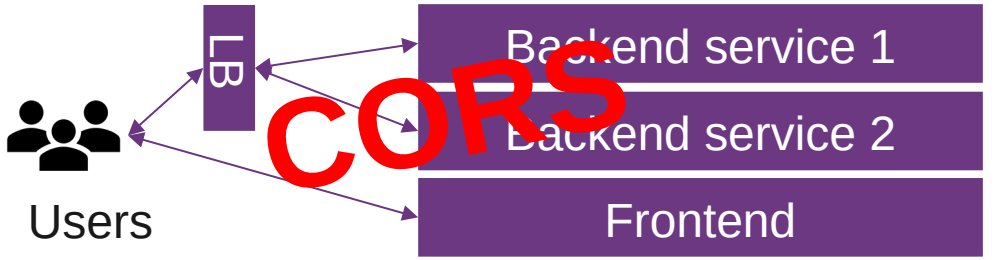
- CORS = Cross-Origin Resource Sharing
  - For security reasons, browsers restrict cross-origin HTTP requests initiated from scripts (among others)
  - Mechanism to instruct browsers that runs a resource from origin A to run resources from origin B
- Solution
  - Use reverse proxy with builtin webserver, e.g., nginx, or user reverse proxy with external webserver.

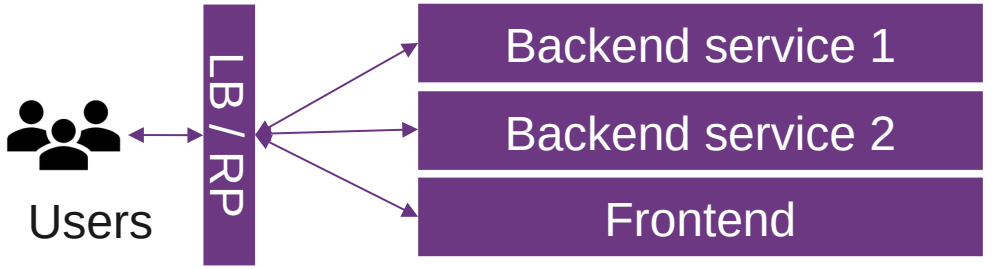
→ The client only sees the same origin for the API and the frontend assets

  - Access-Control-Allow-Origin: https://foo.example

→ For dev: Access-Control-Allow-Origin: *

  - w.Header().Set("Access-Control-Allow-Origin", "*")
- Reverse proxy





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Lecture 7

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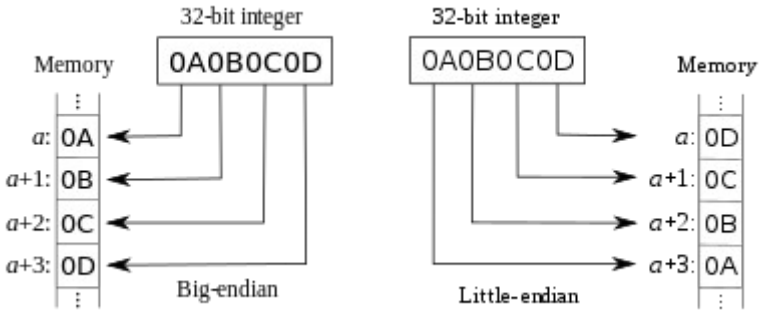
Protocols

- Custom encoding/decoding
  - You control every aspect
  - You send more time on it

```
115     public static boolean decodeHeader(final ByteBuf buffer, final InetSocketAddress recipientSocket,
116             final InetSocketAddress senderSocket, final Message message) {
117         LOG.debug("Decode message. Recipient: {}, Sender:{}.", recipientSocket, senderSocket);
118         final int versionAndType = buffer.readInt();
119         message.version(versionAndType >>> 4);
120         message.type(Type.values()[(versionAndType & Utils.MASK_0F)]);
121         message.protocolType(ProtocolType.values()[versionAndType >>> 30]);
122         message.messageId(buffer.readInt());
123         final int command = buffer.readUnsignedByte();
124         message.command((byte) command);
125         final Number160 recipientID = Number160.decode(buffer);
126
127         //we only get the id for the recipient, the rest we already know
128         final PeerAddress recipient = PeerAddress.builder().peerId(recipientID).build();
129         message.recipient(recipient);
130
131
132         final int contentTypes = buffer.readInt();
133         message.hasContent(contentTypes != 0);
134         message.contentTypes(decodeContentTypes(contentTypes, message));
```

- Little-endian / Big-endian
  - sequential order where bytes are converted into numbers
- Networking, e.g. TCP headers: Big-endian
- Most CPUs e.g., x86: Little-endian, RISC-V: Bi-endianness



[source]

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JSON example

- JSON + REST/HTTP
  - Human-readable text to transmit data
  - Often used for web apps
- 187 bytes

```
func main() {
  fmt.Println("Connecting...")
  req, _ := http.NewRequest("POST", "http://localhost:7000",
    strings.NewReader(`{"code": 5,"message": "Anybody there?"}`))
  req.Header.Set("Content-Type", "application/json")
  client := &http.Client{}
  resp, err := client.Do(req)
  if err != nil {
    panic(err)
  }
  defer resp.Body.Close()
  fmt.Printf("wrote request")
}
```

- Parsing overhead, JSON slower than binary protocol - benchmarks

```
[
    {
        "id": "bitcoin",
        "name": "Bitcoin",
        "symbol": "BTC",
        "rank": "1",
        "price_usd": "9324.08",
        "price_btc": "1.0",
        "24h_volume_usd": "9039300000.0",
        "market_cap_usd": "158560288125",
        "available_supply": "17005462.0",
        "total_supply": "17005462.0",
        "max_supply": "21000000.0",
        "percent_change_1h": "0.46",
        "percent_change_24h": "-0.27",
        "percent_change_7d": "4.5",
        "last_updated": "1525011874“
    }, ...
]
```


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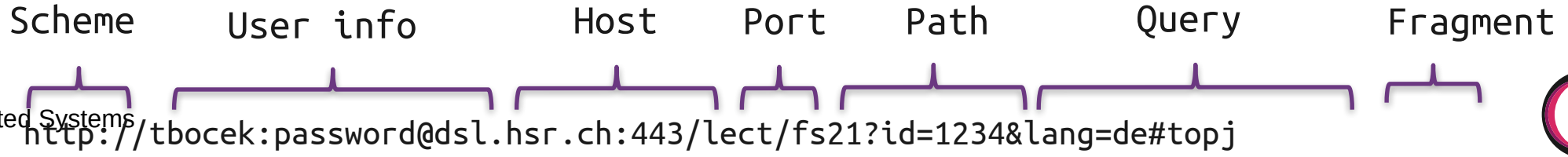
Application Protocol: HTTP

- HTTP (HyperText Transfer Protocol): foundation of data communication for www
- Started in 1989 by Tim Berners-Lee
  - HTTP/1.1 published in 1997
  - HTTP/2 published in 2015
    - More efficient, header compression, multiplexing
  - HTTP/3 wip (April 2022: HTTP/3 protocol is an Internet Draft – not yet final)
- Request / response (resource)
- HTTP resources identified by URL
  - https://dsl.hsr.ch/design/hsr_logo.svg



- Text-based protocol

```
openssl s_client -connect dsl.hsr.ch:443
… TLS handshake …
GET /
```

- Browser sends a bit more…

Request Headers (359 B)

```
Host: dsl.hsr.ch
User-Agent: Mozilla/5.0 (X11; Linux x86_64; rv:73.0) Gecko/20100101 Firefox/73.0
Accept: text/html,application/xhtml+xml,application/xml;q=0.9,image/webp,*/*;q=0.8
Accept-Language: en-US,en;q=0.5
Accept-Encoding: gzip, deflate, br
DNT: 1
Connection: keep-alive
Upgrade-Insecure-Requests: 1
Cache-Control: max-age=0
TE: Trailers
```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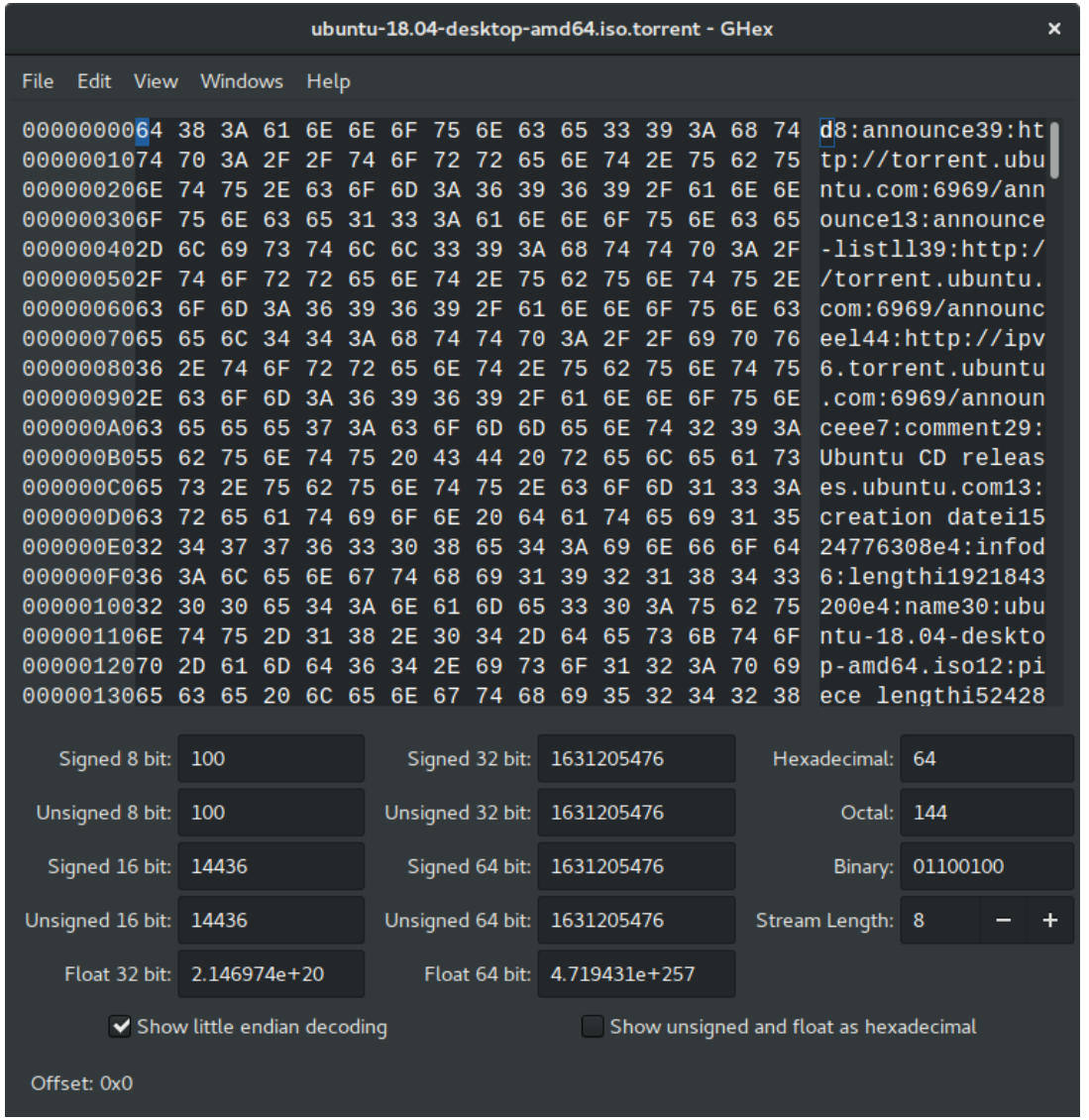
Protocols Bencoding and Others

- Benconding
  - Integers: i42e, Byte string: 4:test, list: l4:testi42ee
  - Map/dictionary: d4:test3:hsr3:tomi42ee
- Use: BitTorrent
- UBJSON
- Cap'n Proto , FlatBuffers
  - Do not serialize, just copy, little-endian
- Apache Arrow
  - Do not serialize, copy, and optimally layout for memory access
- … and many others
- Benchmarks, benchmarks, …

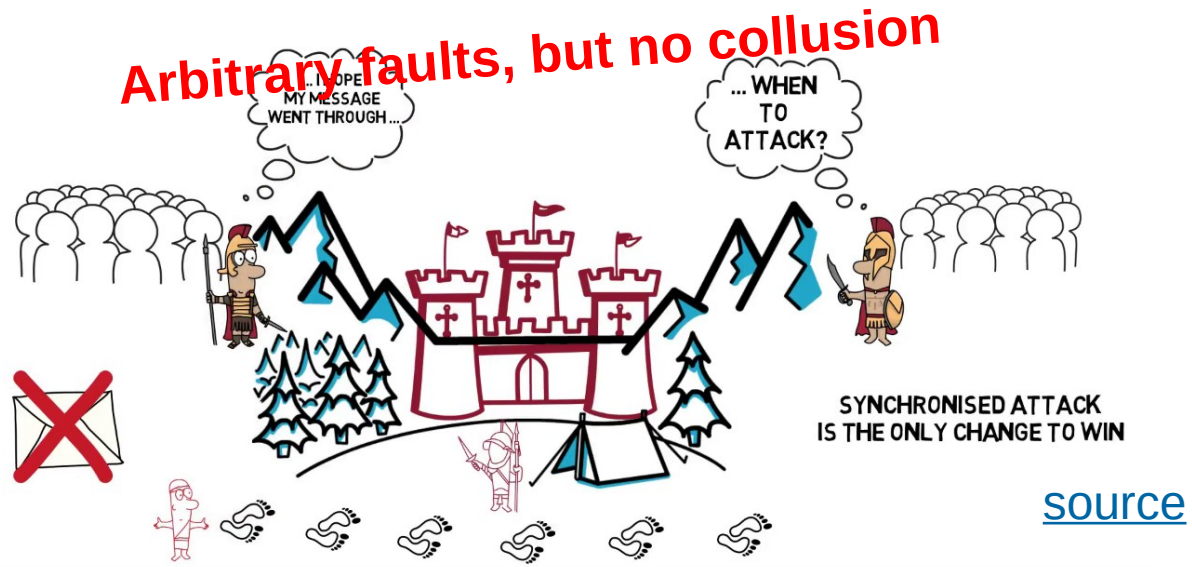


Lecture 8

Consensus

- Definition: Consensus decision-making is a group decision-making process in which group members develop, and agree to **support a decision in the best interest of the whole**.
- A Byzantine fault is an arbitrary fault that occurs during the execution of an algorithm by a distributed system
  - Not only crash, but lie or even collude to reach an advantage
- **“Controlled” Distributed Systems:** your own nodes, your control, no collusion
- Find consensus
  - Paxos, Raft, vDHT, Zookeeper

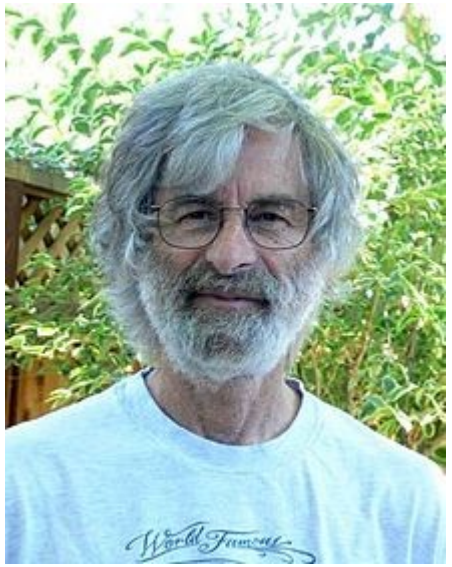
- Often: consensus defines leader
  - Leader creates block
  - Leader adds data
  - Leader creates version
- How to find a leader?



source

Paxos History

- Leslie Lamport discovered the Paxos algorithm in late 1980s
  - Attempt to prove that there was no such algorithm which can tolerate the failure of any number of its processes
  - Until he realized that he created working protocol

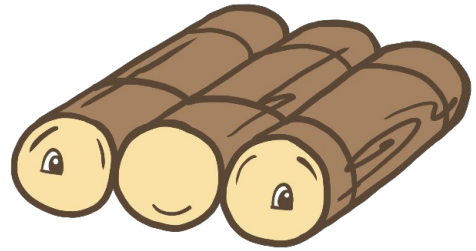


source

- Wrote paper and submitted it to Transactions on Computer Systems (TOCS) in 1990
  - Reviewer: was mildly interesting, but needs significant improvement
  - Leslie Lamport: “so I did nothing with the paper”
- People started to using Paxos to solve problems in distributed systems
- Resubmitted in 1998 to TOCS
  - Accepted without any major changes
- Paxos paper won an ACM SIGOPS Hall of Fame Award in 2012
- Received Turing award in 2013 , also due to Paxos
  - “Turing Award is generally recognized as the highest distinction in computer science and the "Nobel Prize of computing”“ [link]

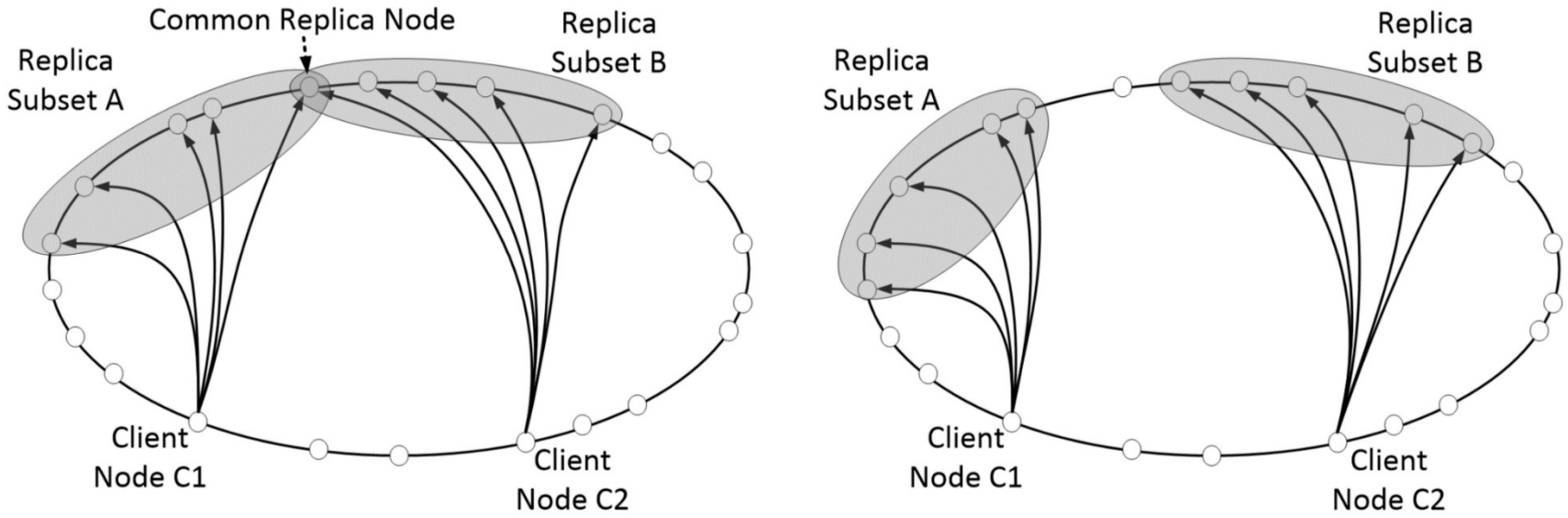
Raft (multi paxos)

- “this makes Raft more understandable than Paxos and also provides a better foundation for building practical systems.” [link]
- RAFT: Reliable, Replicated, Redundant, And Fault-Tolerant
- Follower, Candidate, Leader [link]
  - Raft implements leadership election,
  - Once a leader has been elected, all decision-making within the protocol will then be driven only by the leader
  - Only one leader can exist at a single time
- Each follower has a timeout (typically between 150 and 300 ms) in which it expects the heartbeat from the leader.
  - The system is only available when a leader has been elected and is alive
  - Otherwise, a new leader will be elected and the system will remain unavailable for the duration of the vote
  - Starts election by increasing term counter, voting for itself, and sending a message to all other servers requesting their vote
  - If a higher term is received, become follower, if not, leader



Consistency

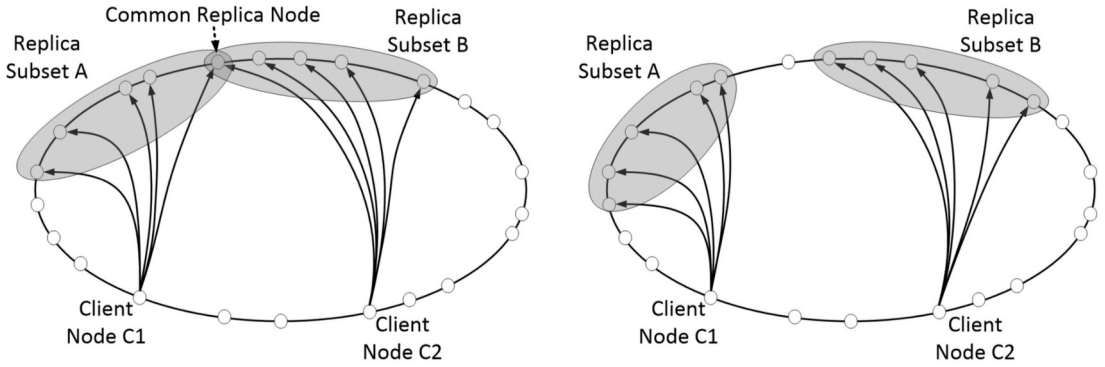
- Consistency in DHTs – vDHT, similarities to Paxos
  - Number = versions, for doing updates
  - Simplified roles (peer)
  - No leader election, works well with churn (not heavy churn)
- CoW, software transactional memory (STM) → for consistent updates. Works for light churn



Lecture 9

CRDT

- (Paxos, why take over larger number?)
  - “acceptors made a promise that no other proposal with a smaller number can make it to consensus” → If acceptor accepted, but its not majority → could stall forever, thus take over large number (link, link)
- L08S10: vDHT
  - A way how to bring consistency to DHTs
  - ~CRDT (operation-based CRDTs)
  - Conflict-free replicated data type (CRDT)
  - ~git but with no merge conflicts
- CRDT must be
  - Commutative x ● y = y ● x
  - Associative (x ● y) ● z = x ● (y ● z)
  - Idempotent x ● x = x



- CRDT Counter (G-Counter)
  - For each machine 1 array position for counter

A	B	C
A:6 B:0 C:0	A:0 B:3 C:0	A:0 B:0 C:9

  - Merge: max of each counter (A:6 B:3 C:9)
    - Old data A:6 B:2 C:9 merge-max / A:5 B:3 C:2 A:6 B:3 C:9
    - Commutative, associative, idempotent

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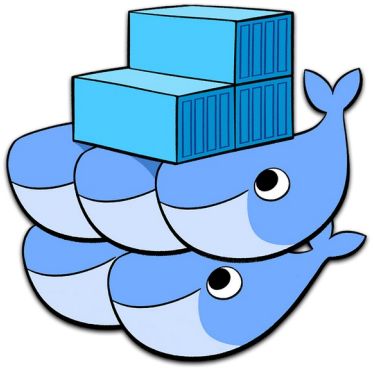
Docker Swarm

- Use docker --context to run/maintain containers on other machines
  - Does not work for docker-compose, could be used with Ansible… “Ansible is also great for bootstrapping Docker itself” [source]
- Docker Swarm
  - Deploy with docker-compose.yml (deploy:)
  - Built into docker
    - docker swarm – manage swarm
    - docker node – manage nodes
  - Scheduler is responsible for placement of containers to nodes
  - Can use the same files, easy to setup?
    - Azure, Google cloud, Amazon



- Kubernetes vs. Docker Swarm
- “Docker Swarm has already lost the battle against Kubernetes for supremacy in the container orchestration space” [link]
- “Kubernetes supports higher demands with more complexity while Docker Swarm offers a simple solution that is quick to get started with.” [link]

Kubernetes

- Kubernetes, K8s
  - Container orchestration (docker)
    - Automated deployment, scaling
  - Started by Google, now with CNCF
- Kubernetes-based PaaS
  - Google, Amazon, Azure (book), Digital Ocean, …
    - Difficult pricing schemes
  - 1.0 released in 2015
  - Package manager Helm released in 2016 (convert docker-compose)
- Why Kubernetes?
  - Containers can crash, machine that runs container can crash (e.g., out of memory)
  - Development: run on one machine, deployment how and where to distribute?
  - Kubernetes manages the lifecycle of containers