



**OST**

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# **Blockchain (BlCh)**

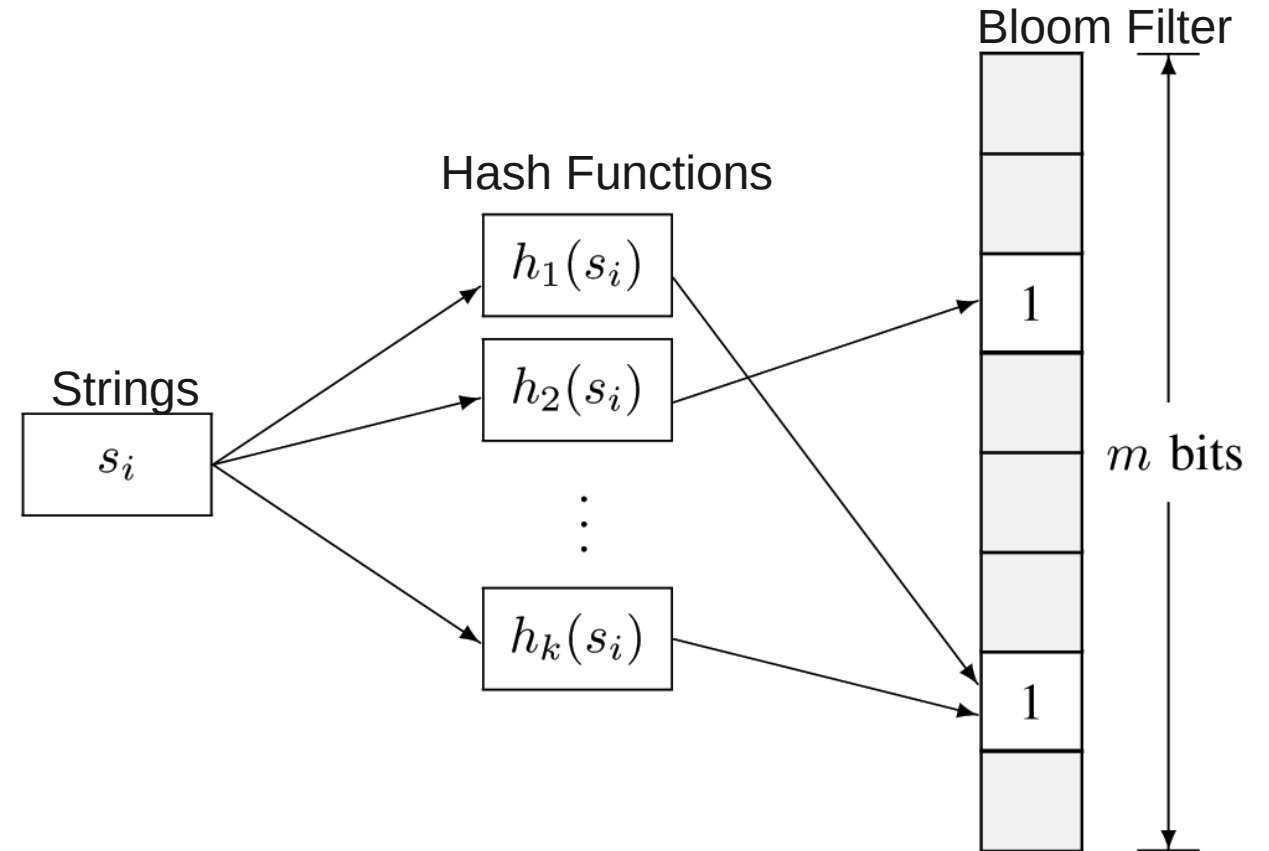
## **Algorithms for P2P Systems**

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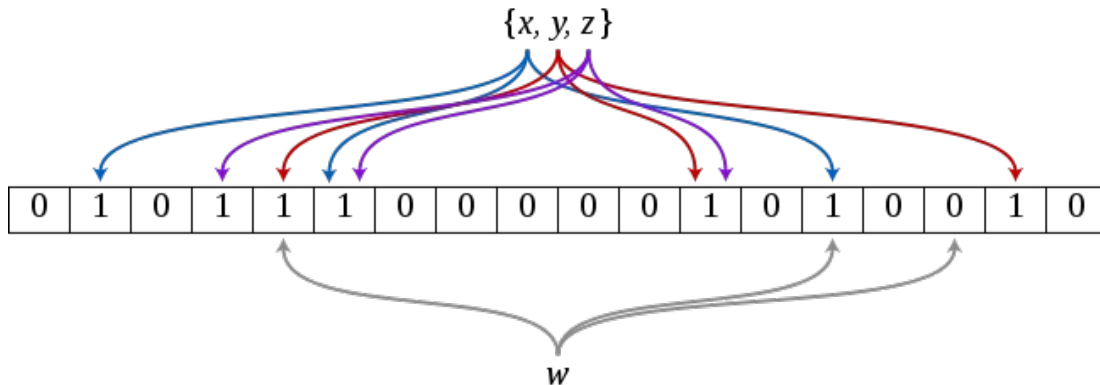
# Bloom Filter

- An array of  $m$  bits, initially all bits set to 0
- A bloom filter uses  $k$  independent hash functions
  - $h_1, h_2, \dots, h_k$  with range  $\{1, \dots, m\}$
- Each input is hashed with every hash function
  - Set the corresponding bits in the vector
- Operations
  - Insertion
    - The bit  $A[h_i(x)]$  for  $1 < i < k$  are set to 1
  - Query
    - Yes if all of the bits  $A[h_i(x)]$  are 1, no otherwise
  - Deletion
    - Removing an element from this simple Bloom filter is impossible



# Query of an Element, $m=18$ , $k=3$

- Insert  $x, y, z$
- Query  $w$



[http://en.wikipedia.org/wiki/Bloom\\_filter](http://en.wikipedia.org/wiki/Bloom_filter)

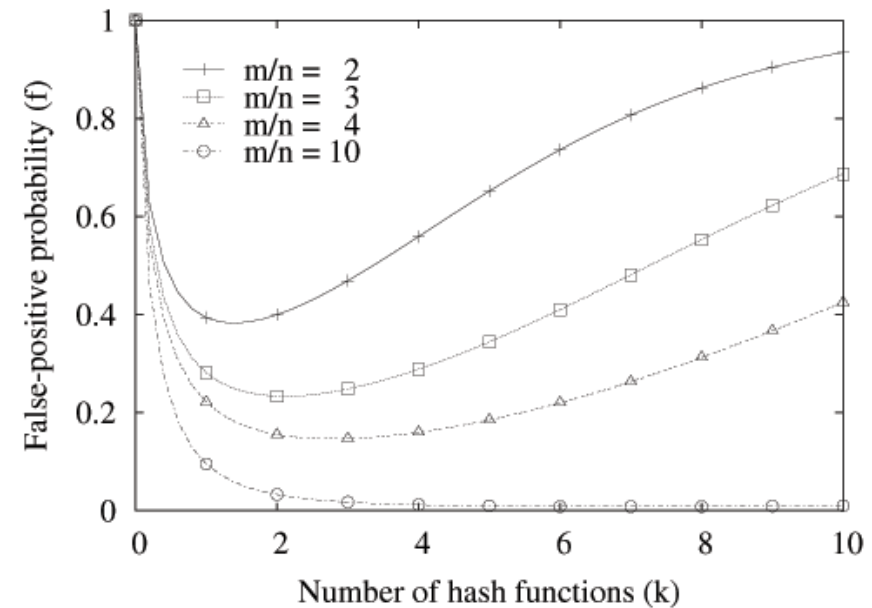
- Example for False-positives
  - Insertions
    - Hash („color printer“) => (1,4,6)
    - Hash („digital camera“) => (3,4,5)
    - Bloom filter (1,3,4,5,6)
  - Query
    - Hash („heat sensor“) => (3,4,6)
    - Matches since bits 3,4,6 are all set to 1
  - Online
- False-negative
  - Query
    - Hash (“color printer“) => (1,4,6) , matches (1,3,4,5,6) → no false-negative

# Properties

- Space Efficiency
  - Any Bloom filter can represent the entire universe of elements
    - In this case, all bits are 1
- No Space Constraints
  - Add never fails
  - But false positive rate increases steadily as elements are added
- Simple Operations
  - Union of Bloom filters: bitwise OR
  - Intersection of Bloom filters: bitwise AND

- No false negative, but false positive
- False-positive probability:
  - $n$  number of strings;  $k$  hash functions;  $m$ -bit vector

$$f = \left(1 - e^{-\frac{nk}{m}}\right)^k$$

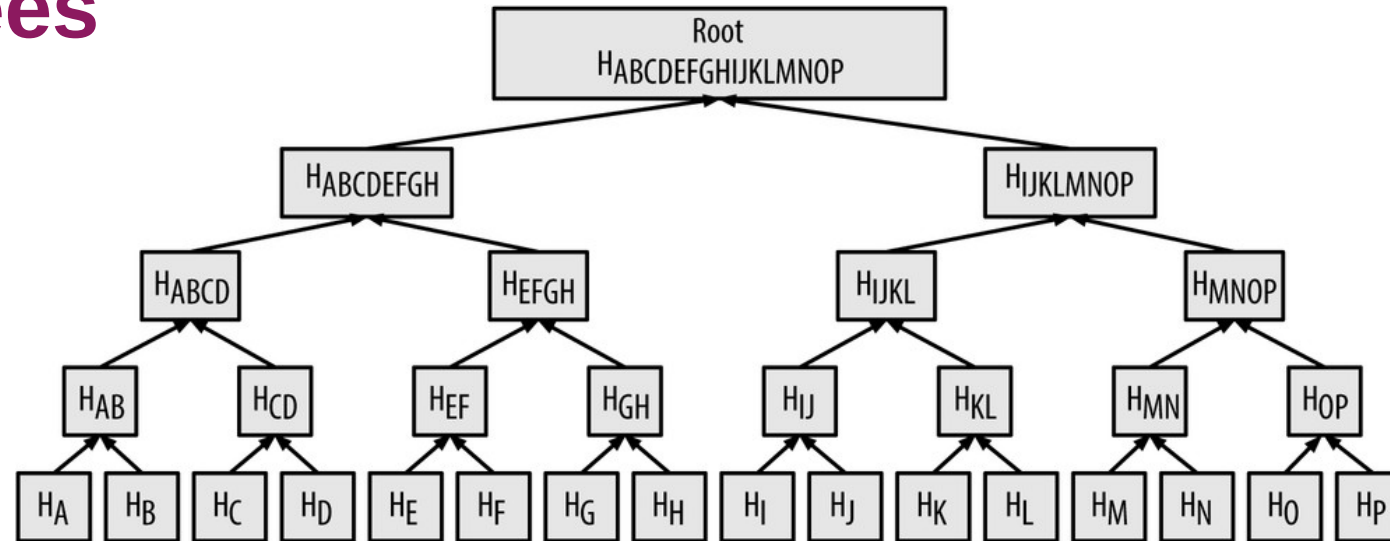


=> Given  $m/n$ , there is an optimal number of hash functions (opt.  $k = m/n \ln 2$ ) (when 50% of the bits are set)

# Bloom Filter Variants

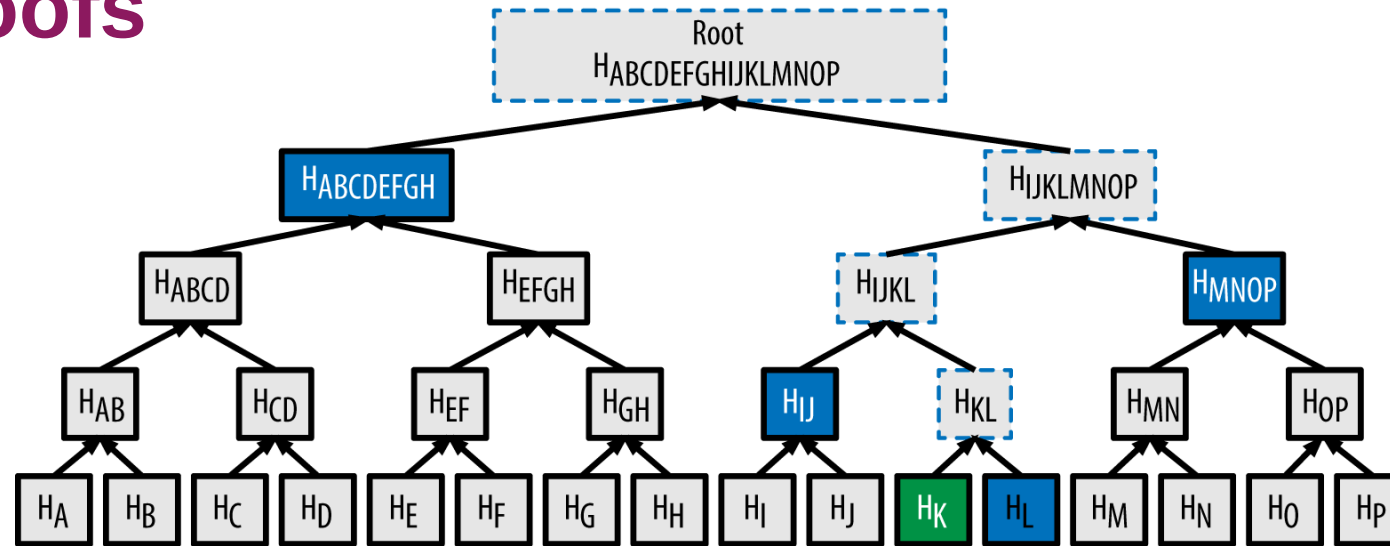
- **Compressed Bloom Filters**
  - When the filter is intended to be passed as a message
  - False-positive rate is optimized for the compressed bloom filter (uncompressed bit vector  $m$  will be larger but sparser)
  - However, compression/decompression, more memory
- **Generalized Bloom Filter**
  - Two type of hash functions  $g_i$  (reset bits to 0) and  $h_j$  (set bits to 1)
  - Start with an arbitrary vector (bits can be either 0 or 1)
  - In case of collisions between  $g_i$  and  $h_j$ , bit is reset to 0
  - Store more info with low false positive
  - Produces either false positives or false negatives
- **Counting Bloom Filters**
  - Entry in the filter not be a single bit but a counter
  - Delete operation possible (decrementing counter)
  - **Variable-Increment Counting Bloom Filter**
- **Scalable Bloom Filter**
  - Adapt dynamically to number of elements, consist of regular Bloom filters
  - “A SBF is made up of a series of one or more (plain) Bloom Filters; when filters get full due to the limit on the fill ratio, a new one is added; querying is made by testing for the presence in each filter”
- Others, e.g., **Cuckoo filter**
- Usage: e.g., **fast search** at LinkedIn

# Merkle Trees



- A Merkle tree is a binary hash tree containing leaf nodes
- Constructed bottom-up, i.e.,
- Used to summarize all transactions in a block
- To prove that a specific transaction is included in a block, a node only needs to produce hashes, constituting a merkle path connecting the specific transaction to the root of the tree.

# Merkle Proofs

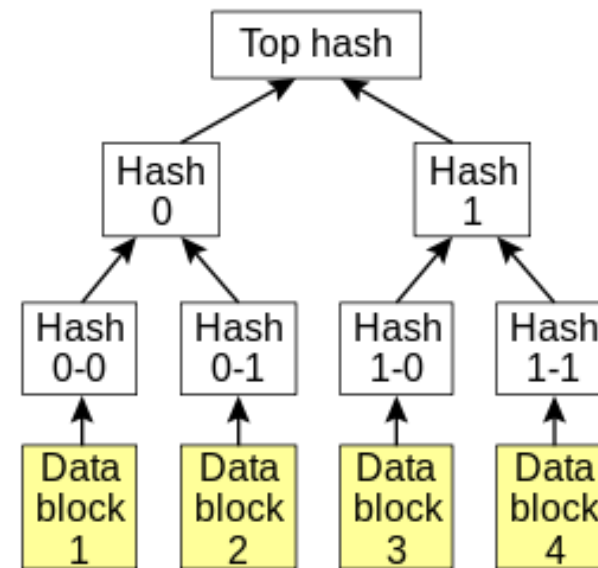


- A node can prove that transaction K is included in the block by producing a merkle path
  - $\log_2 16 = 4$  hashes long

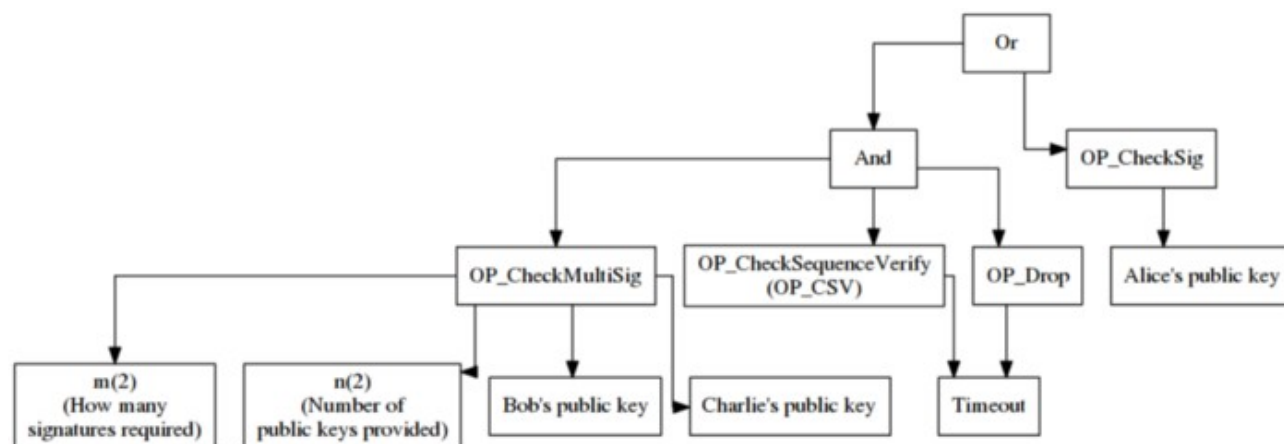


# BitTorrent: Mechanisms

- Magnet links
  - Magnet is URI scheme, does not point to a centralized tracker
    - No centralized tracker: pointer to DHT
    - General purpose, not only for BT
    - magnet:?xl=1000&dn=song1.mp3&xt=urn:tree:tiger:2A3B...
  - tree:tiger → Hash Tree
    - Tree of hashes (|| → concatenation)
    - hash 0 = hash( hash 0-0 || hash 0-1 )
    - hash 1 = hash( hash 1-0 || hash 1-1 )
    - Top hash = hash( hash 0 || hash 1 )
  - Merkle hash / hash tree also seen in Bitcoin blocks (transactions), MAST (Merkalized Abstract Syntax Tree)



[http://en.wikipedia.org/wiki/Hash\\_tree](http://en.wikipedia.org/wiki/Hash_tree)

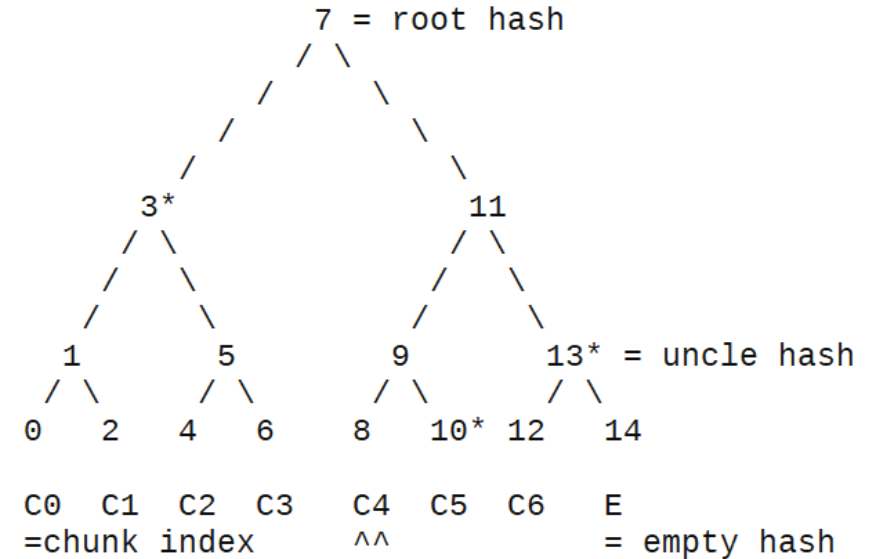


<https://bitcointechtalk.com/what-is-a-bitcoin-merklized-abstract-syntax-tree-mast-33fdf2da5e2f>



# BitTorrent: Mechanisms

- Verification
  - Peer A has top hash (root hash)
  - Peer downloads C4 from peer B
    - create hash 8
  - Need hash 10, 13, 3 (uncle hash)
    - Can be from peer B
  - With 8,10,13,3 can create root hash
    - verify this root hash
- Usage: Blockchain, P2P filesharing, git, Amazons Dynamo, ZFS



The Merkle hash tree of an interval of width  $w=8$

<http://datatracker.ietf.org/doc/draft-ietf-ppsp-peer-protocol/> Section 5.2