

Distributed Systems (DSy)

Protocols

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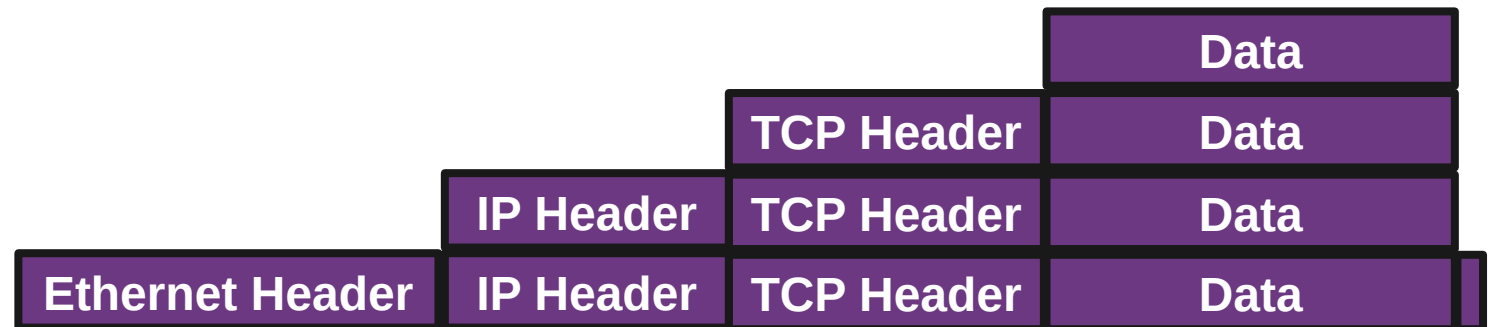
Learning Goals

- Lecture 9 (Protocols)
 - How do network layers work?
 - What are the TCP mechanisms?
 - What are the problems of TCP, and how other protocols (HTTP/3) can improve that

Networking: Layers

- Networking: Each vendor had its own proprietary solution - not compatible with another solution
 - IPX/SPX – 1983, AppleTalk 1985, DECnet 1975, XNS 1977
- Nowadays most vendors build compatible networks hardware/software
 - Cisco, Dell, HP, Huawei, Juniper, Lenovo, Netgear, MicroTik, Ubiquiti, etc.
- Goal of layers: interoperability
 - 1984: ISO 7498 - The Basic Reference Model for Open Systems Interconnection

OSI model	"Internet model"
Application	Application
Presentation	
Session	
Transport	Transport
Network	Internet
Data link	Link
Physical	



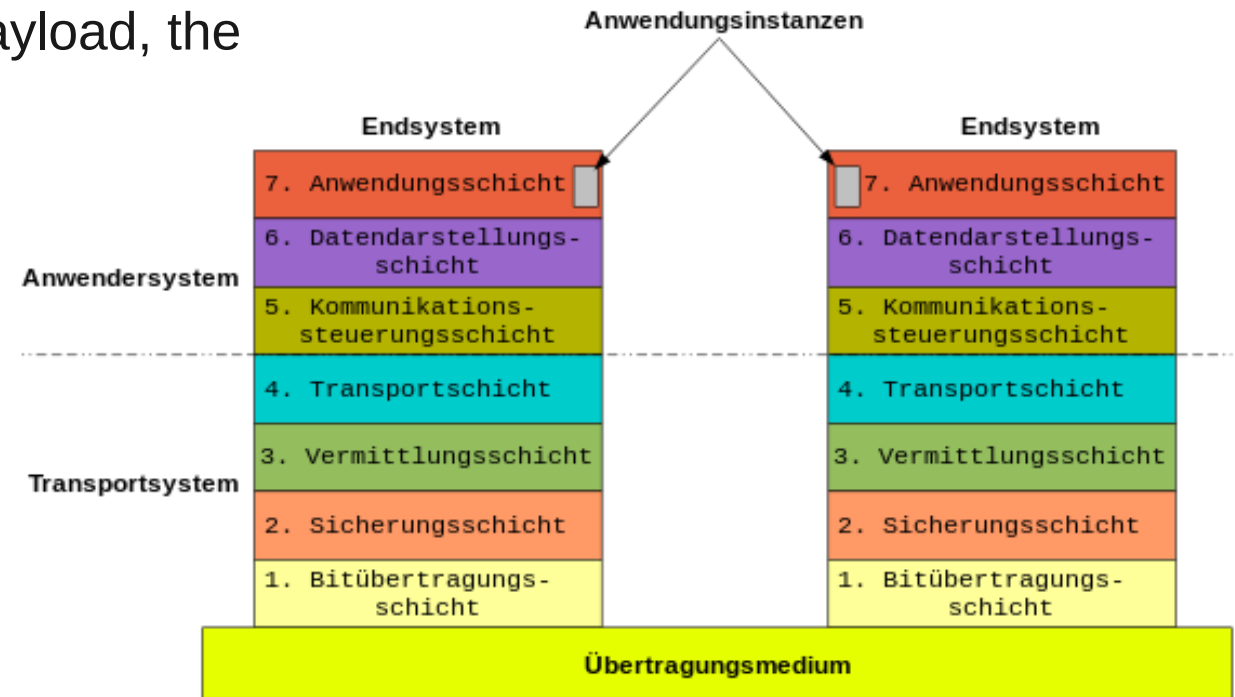
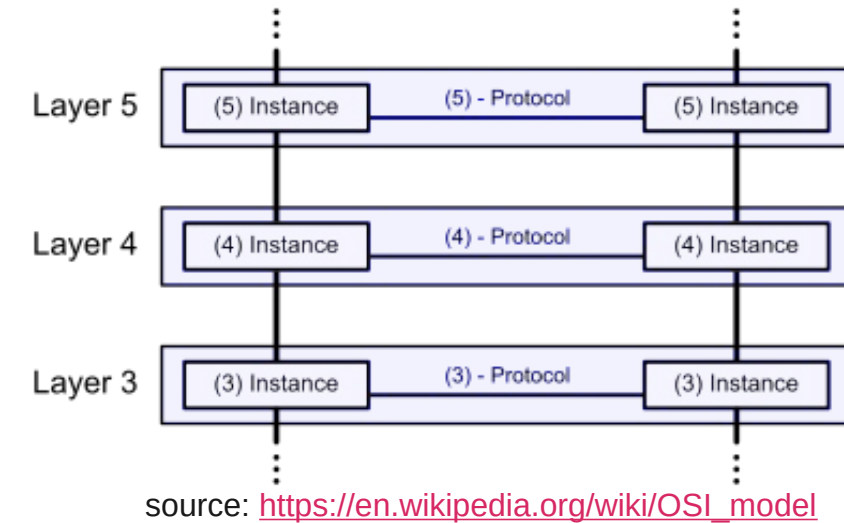
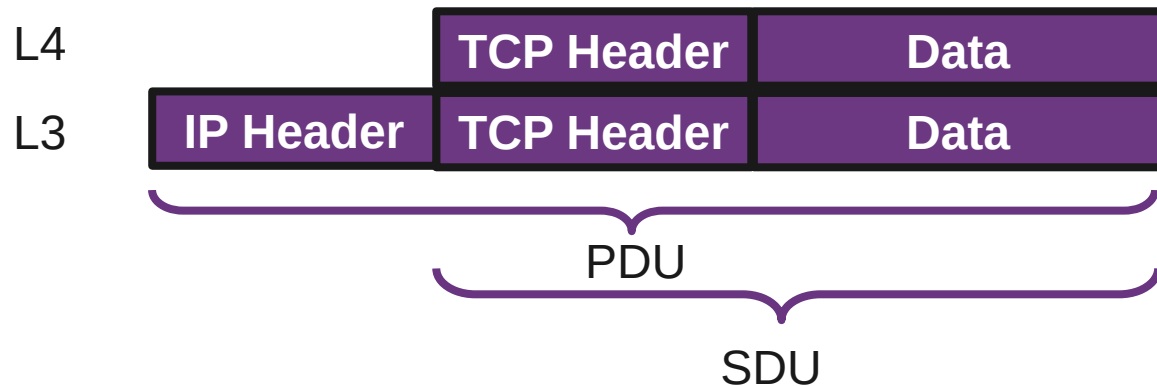
Networking: Definitions

RFC 1122, Internet STD 3 (1989)
Four layers
"Internet model"
Application
Transport
Internet
Link

OSI model
Seven layers
OSI model
Application
Presentation
Session
Transport
Network
Data link
Physical

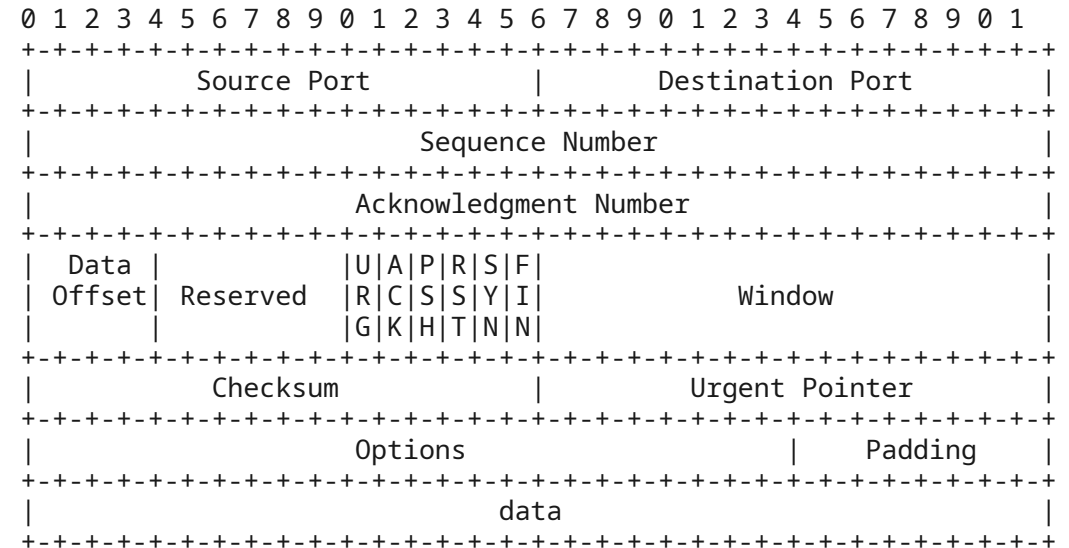
Layer Abstraction

- Protocols enable an entity/instance to interact with an entity/instance at the same layer in another host
- Service definitions: provide functionality to an (N)-layer by an (N-1) layer
- Each **PDU** contains a protocol header and payload, the service data unit (**SDU**). E.g. PDU of L3:



Layer 4 - Transport

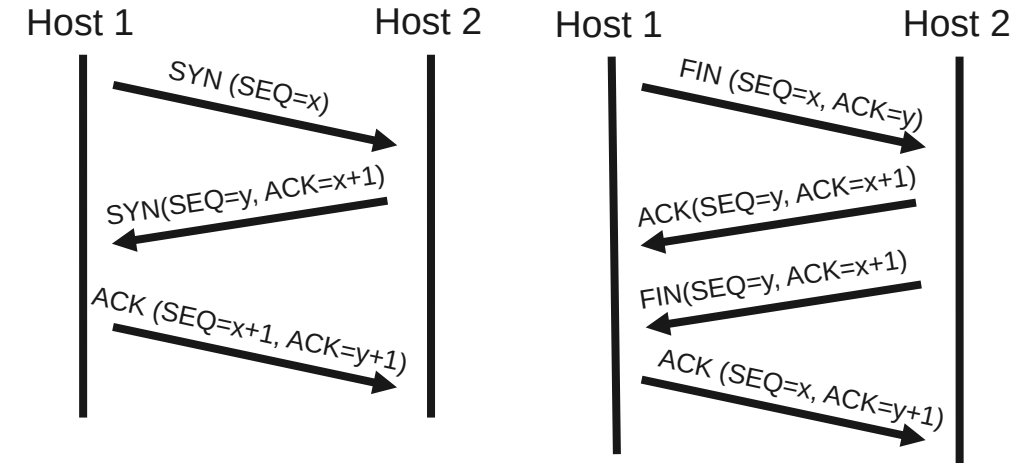
- TCP (Transmission Control Protocol)
 - Reliable (retransmission)
 - Ordered
 - Window – capacity of receiver
 - Checksum – 16bit (crc16)
 - TCP overhead: 20bytes
 - IP overhead: 20bytes
 - Ethernet frame: 18bytes (crc32)
- TCP tries to correct errors; you don't need to worry...
 - Sometimes, you need to worry...



source: <https://datatracker.ietf.org/doc/html/rfc9293>

Layer 4 - TCP

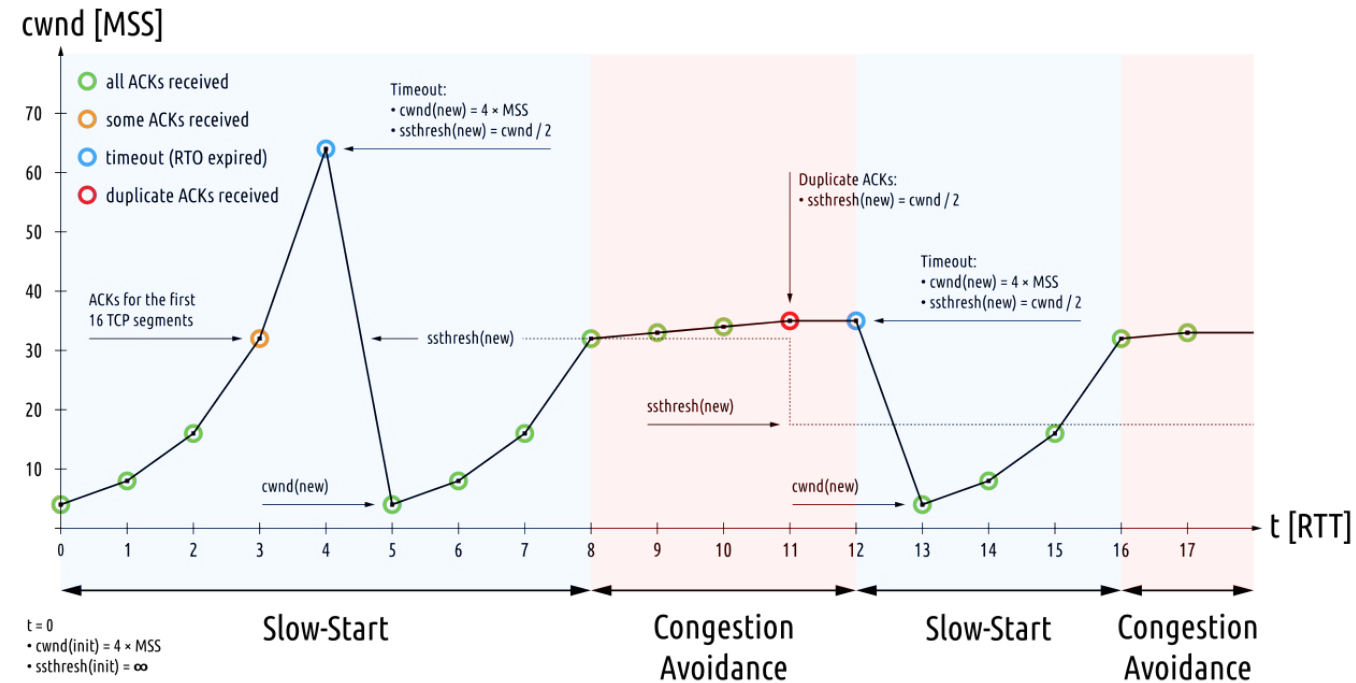
- Connection establishment
 - SYN, SYN-ACK, ACK (three way)
 - Initiates TCP session: initial sequence number is ~random
- Connection termination
 - FIN, ACK + FIN, ACK (three/four way)
 - 3-way handshake, when host 1 sends a FIN and host 2 replies with a FIN & ACK
- Sequences and ACKs
 - Identification each byte of data
 - Order of the bytes → reconstruction
 - Detecting lost data: RTO, DupACK:



- Retransmission timeout
 - If no ACK is received after timeout (e.g. $2 \times \text{RTT}$), resend.
- Duplicate cumulative acknowledgements, selective ACK [[link](#)]
 - ACKs for last consecutive packets
 - 3 times same ACK → retransmit missing packets (fast retransmit)

Layer 4 - TCP

- Flow control
 - Sender is not overwhelming a receiver
 - Back pressure
 - Sliding window:
 - Receiver specifies the amount of additionally received data in bytes that can be buffered
 - Sender up to that amount of data before ACK
- Congestion control
 - slow-start
 - congestion avoidance
- Difference flow/congestion control



source: https://upload.wikimedia.org/wikipedia/commons/thumb/2/24/TCP_Slow-Start_and_Congestion_Avoidance.svg/1280px-TCP_Slow-Start_and_Congestion_Avoidance.svg.png

TCP/IP from an Application Developer View

- Server in go lang ([repo](https://github.com/tbocek/DSy))
 - git clone
<https://github.com/tbocek/DSy>
[\[link\]](#)
 - go run server.go → server
- Listening on TCP port 8081
 - Return string in uppercase
- Node.js version
- Client: nc localhost 8081

```
const net = require('net');
const server = new net.Server();
server.listen(8081, function() {
  console.log('Launching server...');
});

server.on('connection', function(socket) {
  socket.on('data', function(chunk) {
    console.log('Data received from client: $' +
      {chunk.toString()});

    socket.write(chunk.toString().toUpperCase() +
      "\n");
  });
});
```

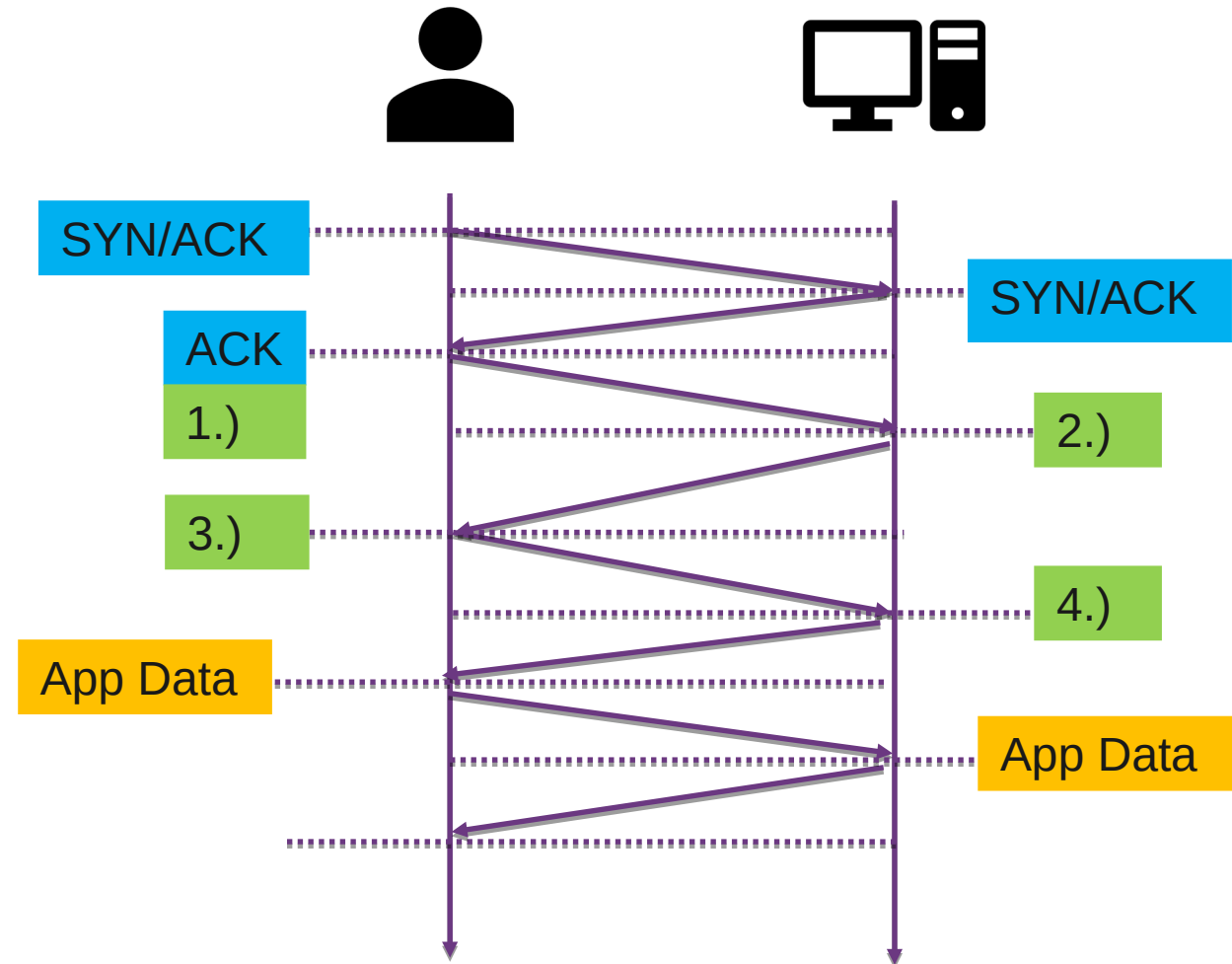
```
package main
import ("bufio"
  "fmt"
  "net"
  "strings")
func main() {
  fmt.Println("Launching server...")
  ln, _ := net.Listen("tcp", ":8081") // listen on all
  interfaces
  for {
    conn, _ := ln.Accept() // accept connection on
    port
    message, _ :=
    bufio.NewReader(conn).ReadString('\n') //read line
    fmt.Print("Message Received:", string(message))
    newMessage := strings.ToUpper(message) //change
    to upper
    conn.Write([]byte(newMessage + "\n")) //send
    upper string back
  }
}
```

TCP Considerations

- Fallacy 2: Latency is zero
 - Nürnberg data center: 12ms, Australia: 300ms
 - Ping ftp.au.debian.org , [Starklink](#) + ~30ms
- Problem: TCP handshake is not flexible
 - You need a handshake (1RT)
 - 1) If you want to make sure the other side accepts packets (and not drop it) - ensure both sides are ready to transmit and receive data
 - 2) If you want to exchange public / private keys
 - TCP supports 1) but not 2)
 - Use another security layer for 2), but a security layer needs at least 1 RT
 - TCP + Security = at least 2 RT
 - Nürnberg + Starlink: $2 \times (12 + 30\text{ms}) = 84\text{ms}$
 - Australian: $(2 \times 300\text{ms}) = 600\text{ms}$
- TCP + Security at least 2 RT
 - DNS query may be required too: 3 RT
 - Old security protocols add RT: 4RT
- Worst case: Starlink/Australia/DNS/TCP/old sec: 1.4s before data can be sent → new protocols on the way (HTTP/3)

Layer 4 – TCP + TLS

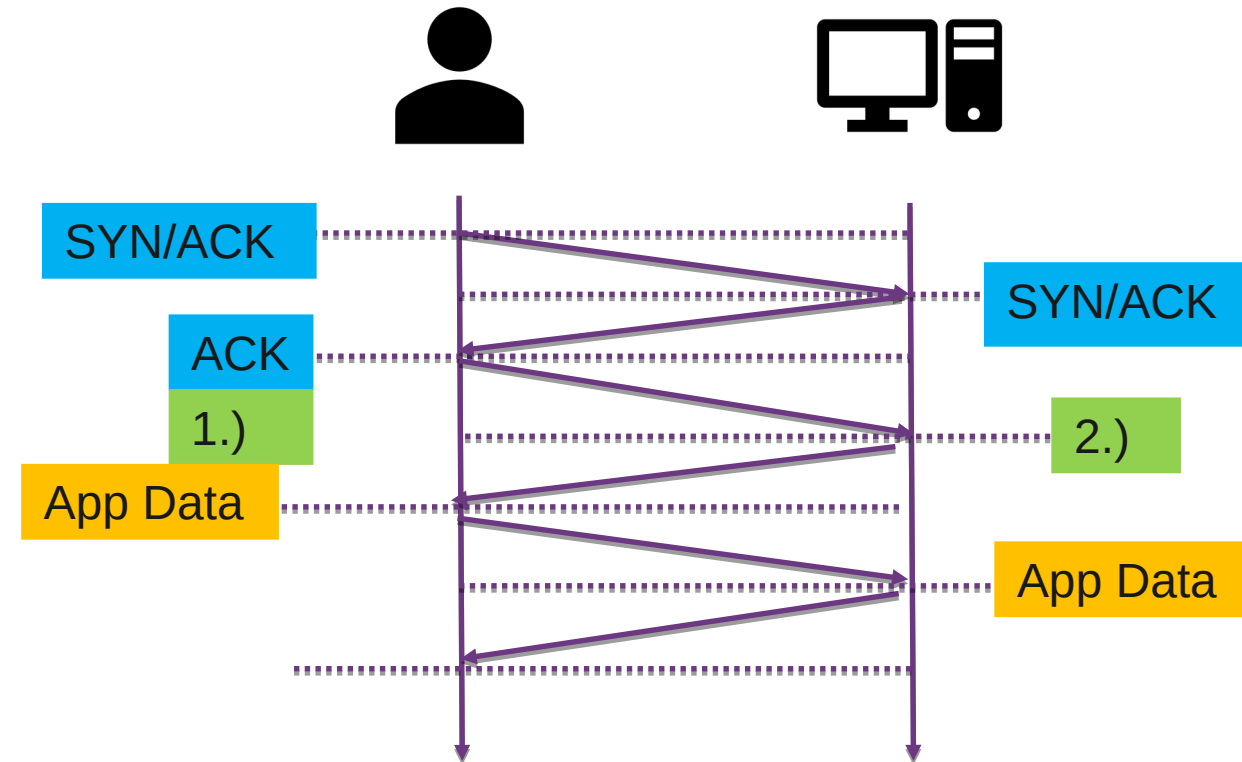
- Security: Transport Layer Security (TLS)
 1. "client hello" lists cryptographic information, TLS version, ciphers/keys
 2. "server hello" chosen cipher, the session ID, random bytes, digital certificate (checked by client), optional: "client certificate request"
 3. Key exchange using random bytes, now server and client can calc secret key
 4. "finished" message, encrypted with the secret key
- 3 RTT to send first byte, 4RTT to receive first byte



```
PING sydney.edu.au (129.78.5.8) 56(84) bytes of data.  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=1 ttl=233 time=307 ms  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=2 ttl=233 time=305 ms  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=3 ttl=233 time=305 ms
```

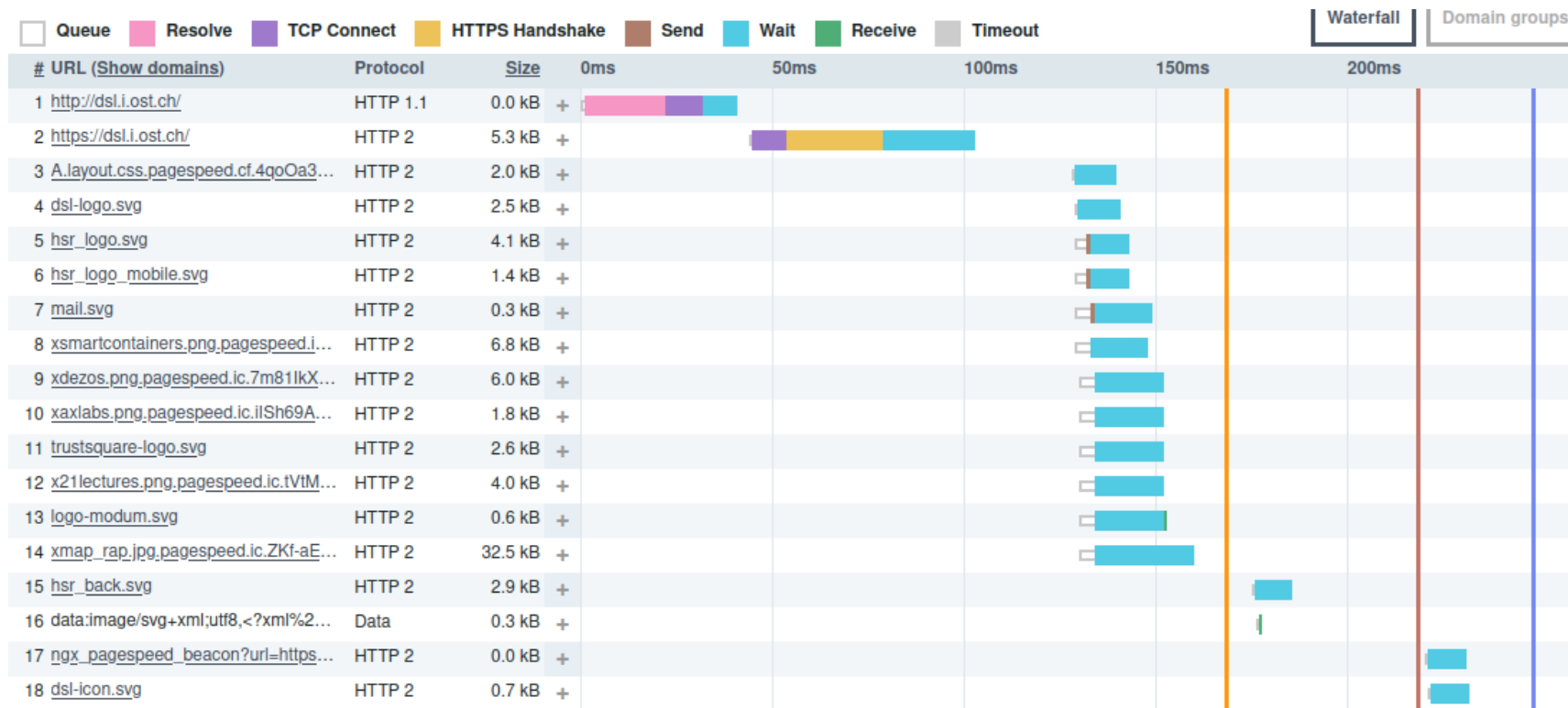
Layer 4 – TCP + TLS

- TCP + TLS handshake:
 - 1RTT - Ping to Australia: 329ms
 - 3RTT = 987ms! No data sent yet
- TLS 1.3, finished Aug 2018
 - 1 RTT instead of 2
 - 1.) Client Hello, Key Share
 - 2.) Server Hello, key Share, Verify Certificate, Finished
 - 0 RTT possible, for previous connections
- 95% of browsers used already support it



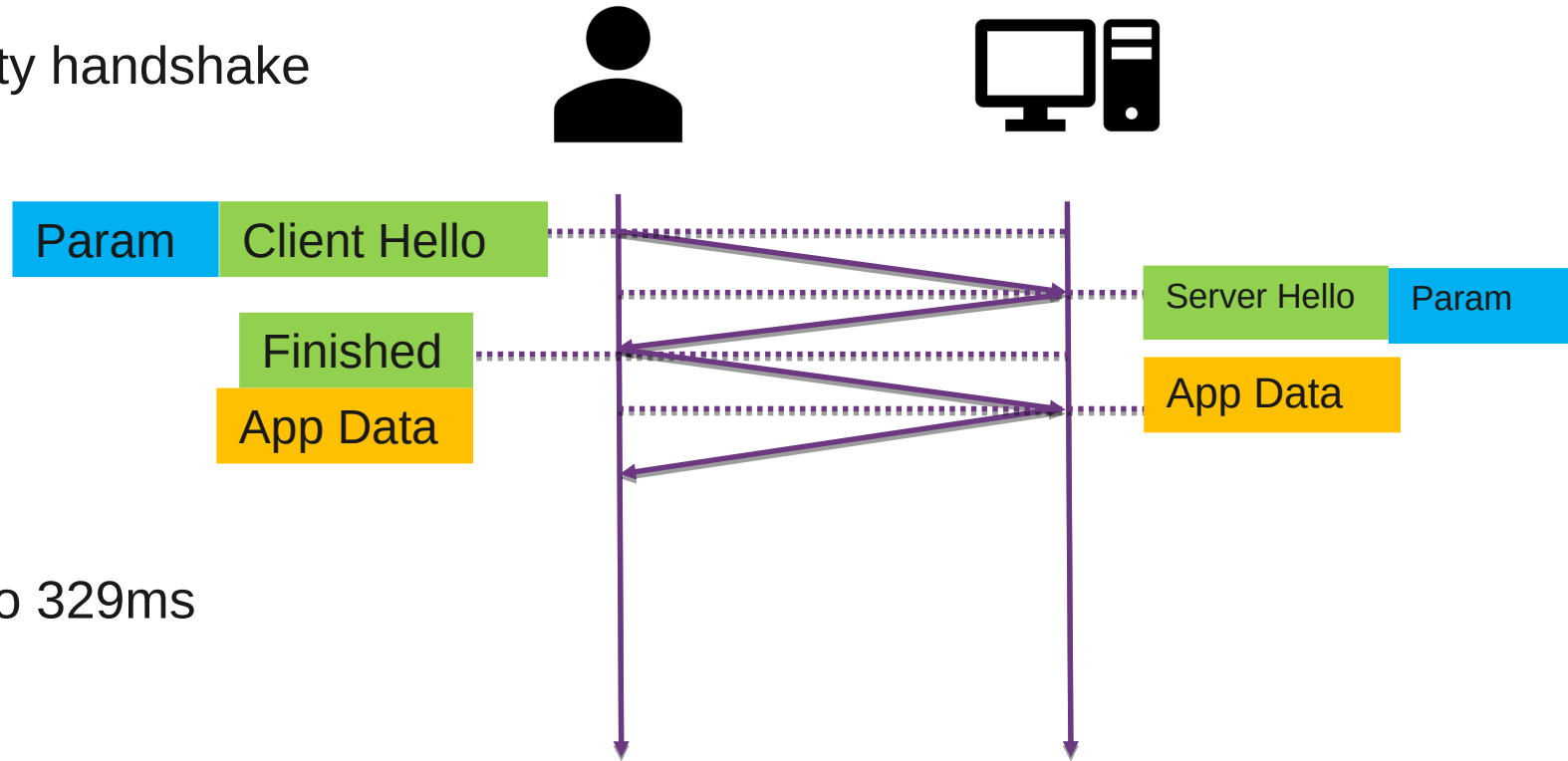
Layer 4 – TCP + TLS

- Website Speed Test [\[link\]](#)
- Resolve → DNS, TCP Connect → TCP Handshake, HTTPS Handshake → TLS/SSL



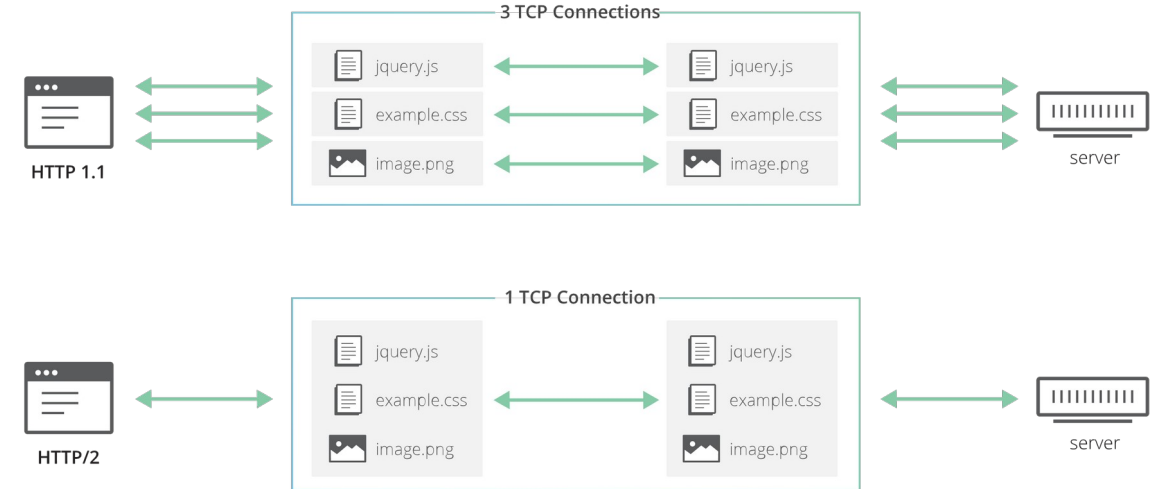
QUIC / HTTP3

- QUIC: 1RTT connection + security handshake
 - For known connections: 0RTT
 - Built in security
 - Reports
 - Facebook
 - state of HTTP
- Example Australia: from 987ms to 329ms



QUIC / HTTP3

- Multiplexing in HTTP/2
 - [HTTP/1 → HTTP/2](#)
- HTTP/2: Head-of-line blocking
 - One packet loss, TCP needs to be ordered
 - QUIC can multiplex requests: one stream does not affect others
- HTTP/3 is great, but...
 - NAT → SYN, ACK, FIN, conntrack knows when connection ends, not with QUIC, timeouts, new entries, many entries
 - HTTP header compression, referencing previous headers
 - Many TCP [optimizations](#)



source: <https://blog.cloudflare.com/the-road-to-quic/>

(#1) GET b.css part 1 ✓

(#2) GET b.js part 1 ✗

(#3) GET b.css part 2 ⌚

(s1) GET b.css part 1 ✓

(s2) GET b.js part 1 ✗

(s1) GET b.css part 2 ✓

Layer 4 - Transport

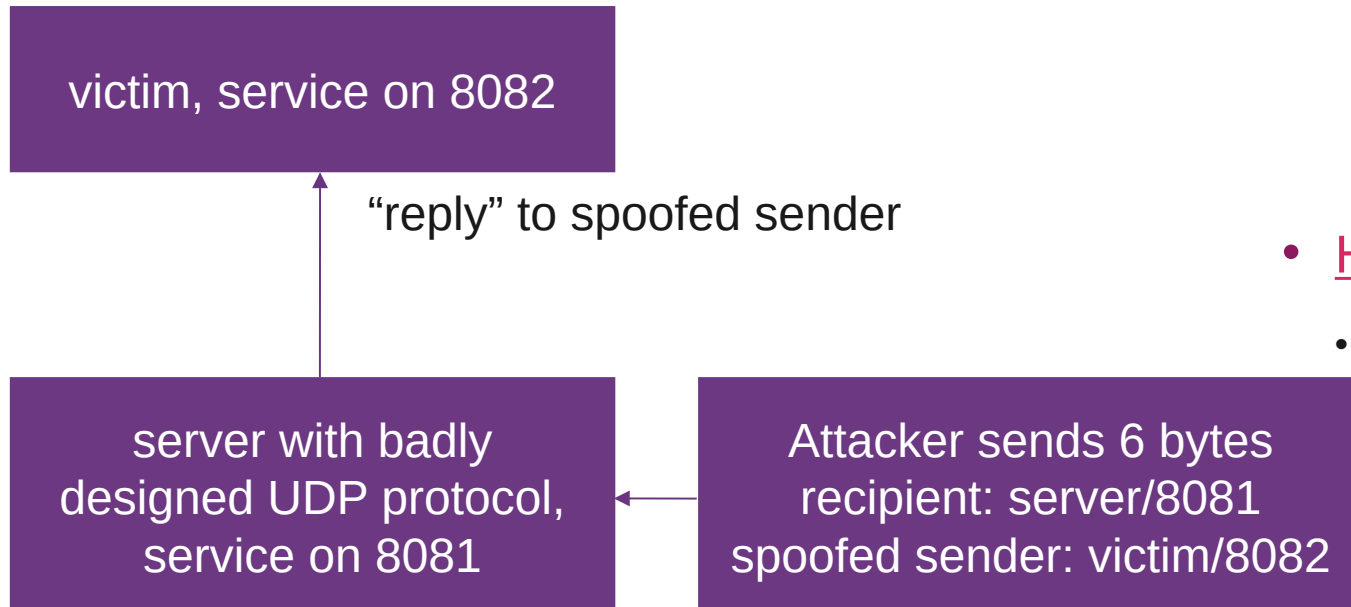
- User Datagram Protocol (UDP)
 - UDP is used for DNS, streaming audio and video
 - Simple connectionless communication model
 - No guarantee
 - Delivery
 - Ordering
 - Duplicate protection



- SCTP (Stream Control Transmission Protocol)
 - Message-based
 - Allows data to be divided into multiple streams
 - Syn cookies - SCTP uses a four-way handshake with a signed cookie.
 - Multi-homing multiple IP addresses of endpoints
 - Not widely used: “We have been deploying SCTP in several applications now, and encountered significant problem with SCTP support in various home routers.”
 - E.g., OpenWRT – not enabled by default
 - E.g., UFW - Uncomplicated Firewall – not supported
 - SCTP used by WebRTC, but tunneled over UDP

Layer 4 - Transport

- DDoS Amplification Attacks
 - Request 10 bytes, reply 100 bytes → factor 10
- Local demo with server-ra/victim, and hping3
 - `hping3 --udp IP -p 8081 -E test.tmp -d 6 -s 8082 -c 1`



- Attacker in go/Java/node/c#
 - You need to spoof UDP packets, typically not supported in those languages
 - Go: `func DialUDP(network string, laddr, raddr *UDPAddr) (*UDPConn, error)`
 - laddr: we need to set here the victims IP/port
 - But go tries to bind to that
 - Not yours: "bind: cannot assign requested address"
- Hping3: Pen test tool
 - hping3 is a command-line oriented IP, TCP, UDP, ICMP and RAW-IP packet assembler

Comparison – Transport Layer

TCP *

- Transport layer
- Connection oriented
- Reliable transfer
- Streams
- Guaranteed order
- Widely used – HTTP/1, HTTP/2
- Flow and congestion control
- Heavyweight
- Error checking and recovery

UDP *

- Transport layer
- Connection less
- Unreliable transfer
- Messages
- Unordered
- Widely used – DNS, HTTP/3
- No flow, congestion
- Lightweight
- Error checking, no recovery

SCTP *

- Transport layer
- Connection oriented
- Reliable transfer
- Messages
- User can choose
- [WebRTC](#)
- Flow and congestion control
- Heavyweight
- Error checking and recovery

QUIC *

- Transport layer*
- Connection oriented
- Reliable transfer
- Multistream
- Guaranteed order
- HTTP/3
- Flow and congestion control
- Heavyweight
- Integrity check