



OST

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Distributed Systems (DSy)

Protocols 2

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Learning Goals

- Lecture 4 (Protocols, part 2)
 - How can new protocols improve latency?
 - What is an amplification attack?

Wireshark – sometimes needed when designing protocols

- Decrypt TLS
- export `SSLKEYLOGFILE=/tmp/keylogfile.txt`

The screenshot displays the Wireshark interface for a network capture titled '*enx106530e2c54d'. The filter bar shows the expression 'ip.src==152.96.80.48 || ip.dst==152.96.80.48'. The packet list pane shows a series of packets, with packet 30 highlighted. The packet details pane for packet 30 shows the following information:

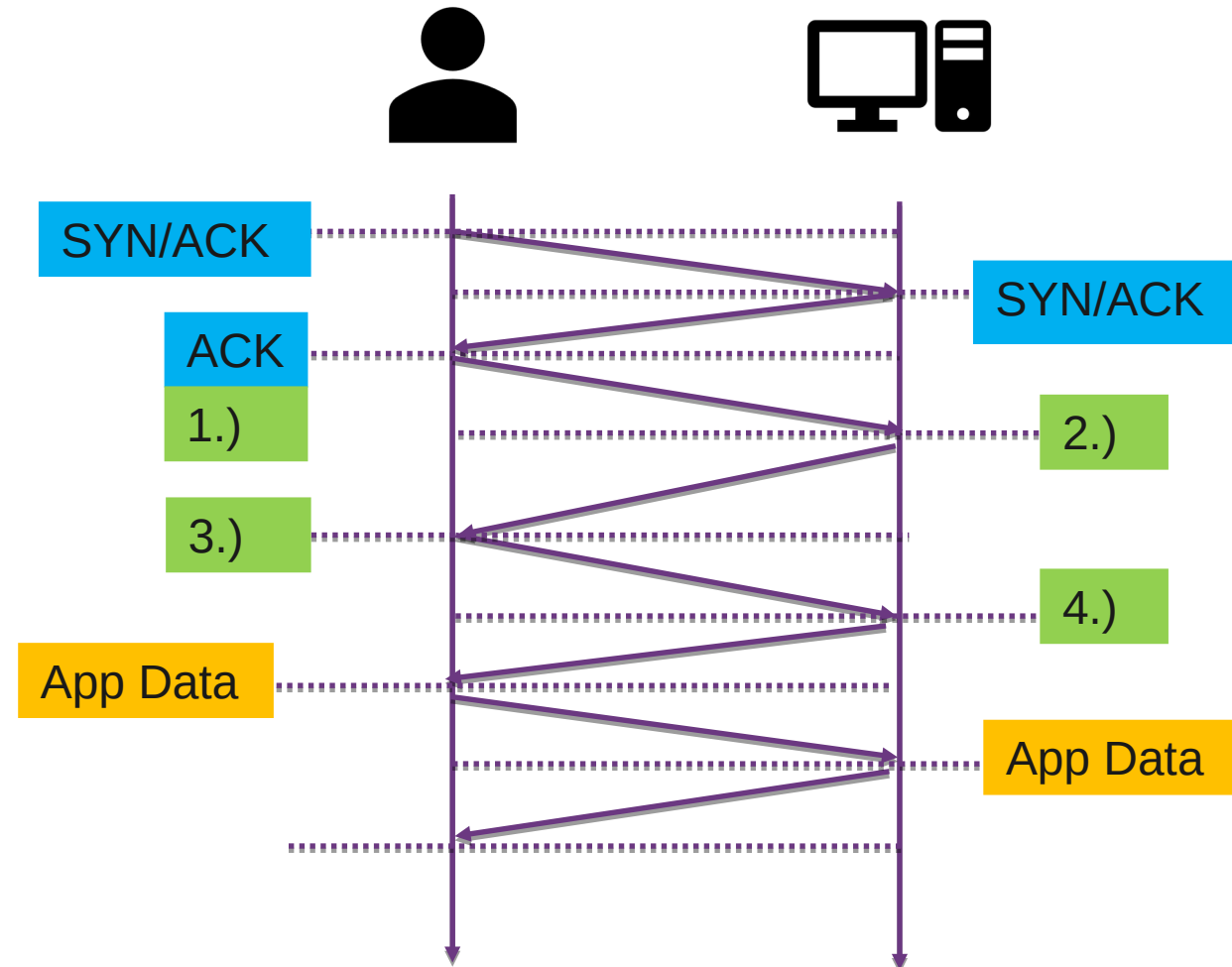
- Frame 30: 74 bytes on wire (592 bits), 74 bytes captured (592 bits) on interface 0
- Ethernet II, Src: Dell_e2:c5:4d (10:65:30:e2:c5:4d), Dst: Cisco_ff:fd:90 (00:08:e3:ff:fd:90)
- Internet Protocol Version 4, Src: 152.96.214.84, Dst: 152.96.80.48
- Transmission Control Protocol, Src Port: 50908, Dst Port: 443, Seq: 0, Len: 0
 - Source Port: 50908
 - Destination Port: 443
 - [Stream index: 5]
 - [TCP Segment Len: 0]
 - Sequence number: 0 (relative sequence number)
 - [Next sequence number: 0 (relative sequence number)]
 - Acknowledgment number: 0
 - 1010 ... = Header Length: 40 bytes (10)
 - Flags: 0x0c2 (SYN, ECN, CWR)
 - 000. = Reserved: Not set

The packet bytes pane shows the raw hex and ASCII data of the SYN packet:

```
0000 00 08 e3 ff fd 90 10 65 30 e2 c5 4d 08 00 45 00 .....e 0...M..E.
0010 00 3c 2f d6 40 00 40 06 b3 a0 98 60 d6 54 98 60 </@.@.....T.
0020 50 30 c6 dc 01 bb 7d 4a 53 f3 00 00 00 00 a0 c2 P0.....}J S.....
0030 72 10 57 74 00 00 02 04 05 b4 04 02 08 0a ba 62 r.Wt.....b
0040 7d 59 00 00 00 00 01 03 03 07 }Y.....
```

Layer 4 – TCP + TLS

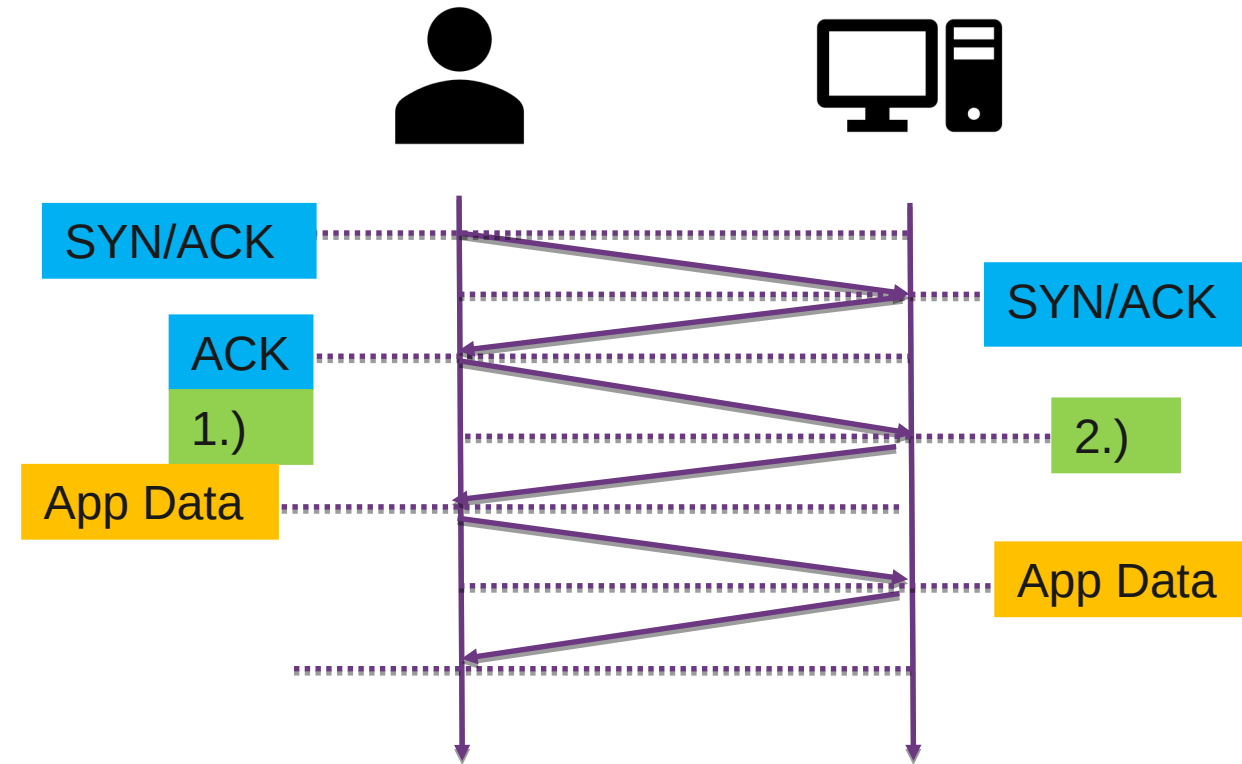
- Security: Transport Layer Security (TLS)
 1. "client hello" lists cryptographic information, TLS version, ciphers/keys
 2. "server hello" chosen cipher, the session ID, random bytes, digital certificate (checked by client), optional: "client certificate request"
 3. Key exchange using random bytes, now server and client can calc secret key
 4. "finished" message, encrypted with the secret key
- 3 RTT to send first byte, 4RTT to receive first byte



```
PING sydney.edu.au (129.78.5.8) 56(84) bytes of data.  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=1 ttl=233 time=307 ms  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=2 ttl=233 time=305 ms  
64 bytes from scilearn.sydney.edu.au (129.78.5.8): icmp_seq=3 ttl=233 time=305 ms
```

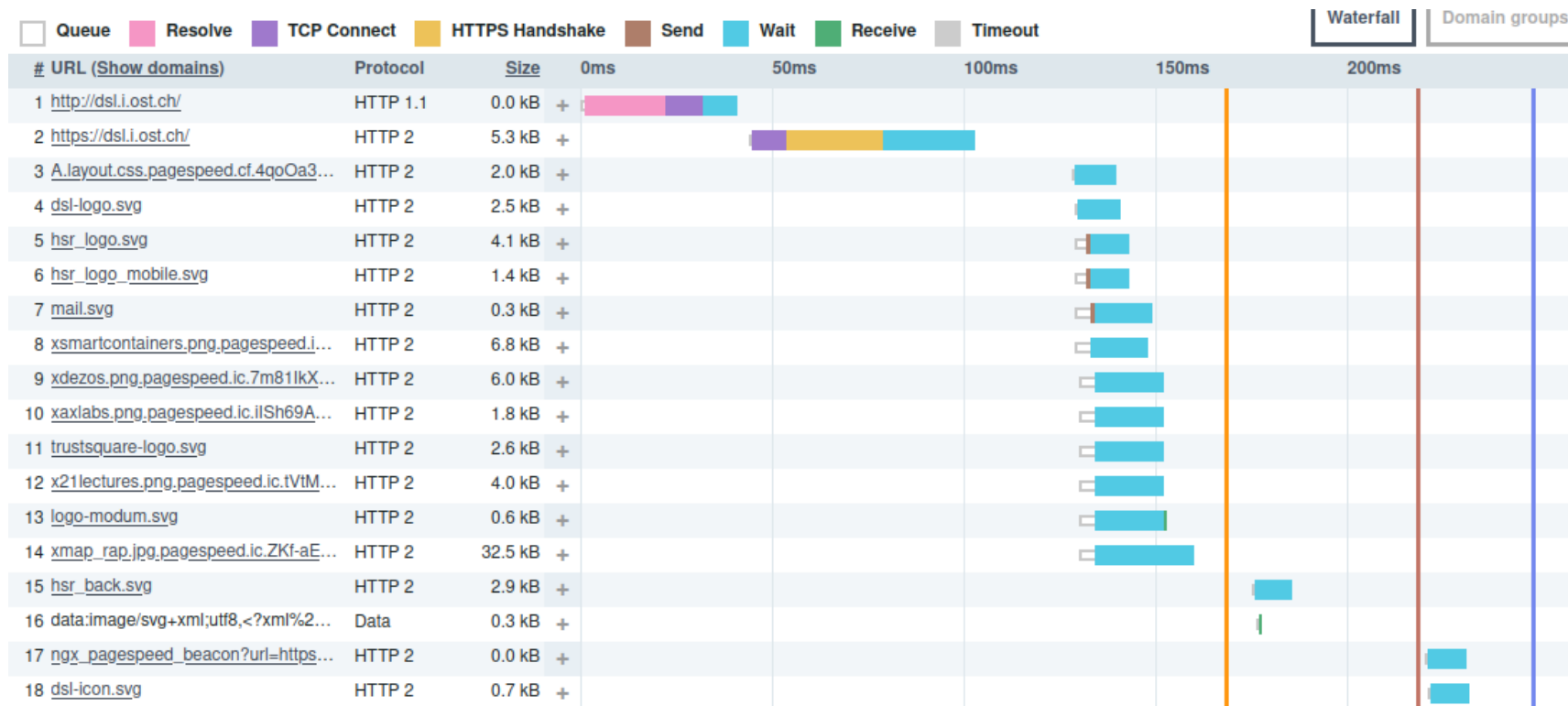
Layer 4 – TCP + TLS

- TCP + TLS handshake:
 - 1RTT - Ping to Australia: 329ms
 - 3RTT = 987ms! No data sent yet
- TLS 1.3, finished Aug 2018
 - 1 RTT instead of 2
 - 1.) Client Hello, Key Share
 - 2.) Server Hello, key Share, Verify Certificate, Finished
 - 0 RTT possible, for previous connections, losing perfect forward secrecy
- 95% of browsers used already support it



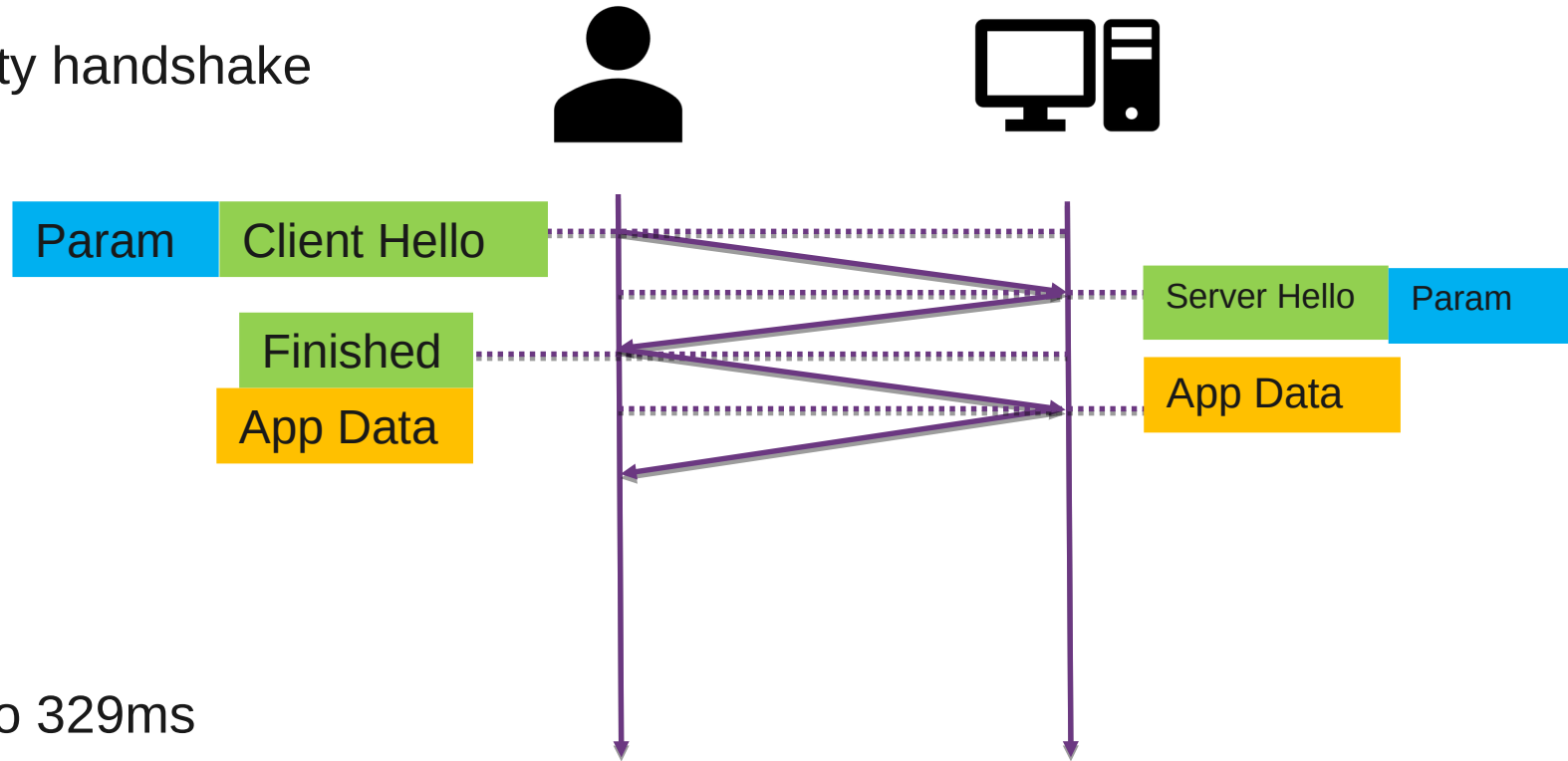
Layer 4 – TCP + TLS

- Website Speed Test [\[link\]](#)
 - Resolve → DNS, TCP Connect → TCP Handshake, HTTPS Handshake → TLS/SSL



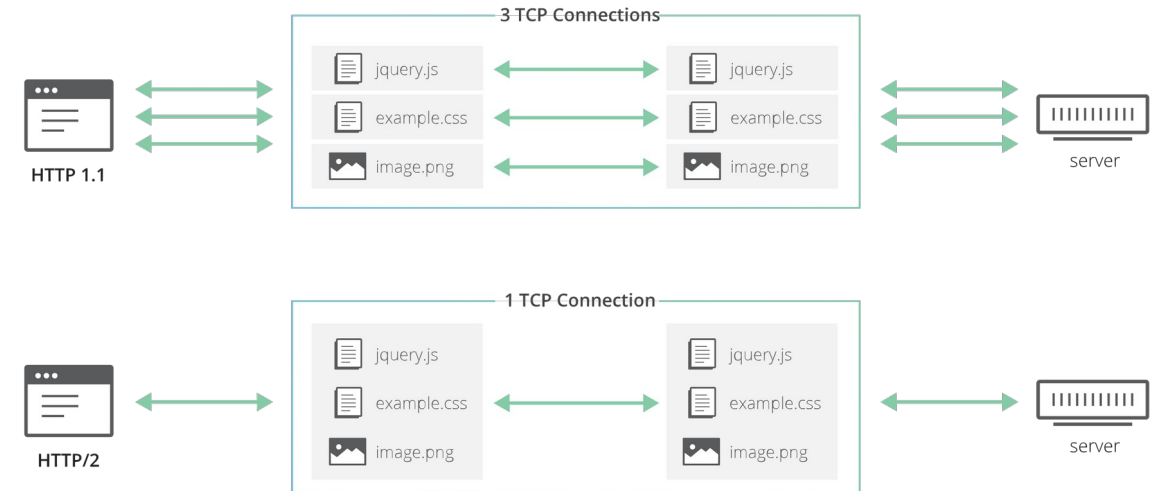
QUIC / HTTP3

- QUIC: 1RTT connection + security handshake
 - For known connections: 0RTT
 - [Built in security](#)
 - “Google's 'QUIC' TCP alternative slow to excite anyone outside Google” [link] ([9%](#), [25%](#), 75%)
 - [Facebook](#)
 - [Cloudflare](#), [state of HTTP](#)
- Example Australia: from 987ms to 329ms



QUIC / HTTP3

- Multiplexing in HTTP/2
 - [HTTP/1 → HTTP/2](#)
- HTTP/2: Head-of-line blocking
 - One packet loss, TCP needs to be ordered
 - QUIC can multiplex requests: one stream does not affect others
- HTTP/3 is great, but...
 - NAT → SYN, ACK, FIN, conntrack knows when connection ends, not with QUIC, timeouts, new entries, many entries
 - HTTP header compression, referencing previous headers
 - Many TCP [optimizations](#)



source: <https://blog.cloudflare.com/the-road-to-quic/>



Layer 4 - Transport

- User Datagram Protocol (UDP)
 - UDP is used for DNS, streaming audio and video
 - Simple connectionless communication model
 - No guarantee
 - Delivery
 - Ordering
 - Duplicate protection



- SCTP (Stream Control Transmission Protocol)
 - Message-based
 - Allows data to be divided into multiple streams
 - Syn cookies - SCTP uses a four-way handshake with a signed cookie.
 - Multi-homing multiple IP addresses of endpoints
 - Not widely used: “
We have been deploying SCTP in several applications now, and encountered significant problem with SCTP support in various home routers.
”
 - E.g., OpenWRT – not enabled by default
 - E.g., UFW - Uncomplicated Firewall – not supported
 - SCTP used by WebRTC, but tunneled over UDP

UDP example

- UDP Server (Java)

```
import java.net.*;

class Server
{
    public static void main(String args[]) throws
Exception
    {
        DatagramSocket serverSocket = new
DatagramSocket(8081);
        byte[] receiveData = new byte[1024];
        while(true)
        {
            DatagramPacket receivePacket = new
DatagramPacket(receiveData, receiveData.length);
            serverSocket.receive(receivePacket);
            String s = new
String( receivePacket.getData());
            System.out.println("Message Received: " +
s);
        }
    }
}
```

- UDP Client (golang)

```
package main

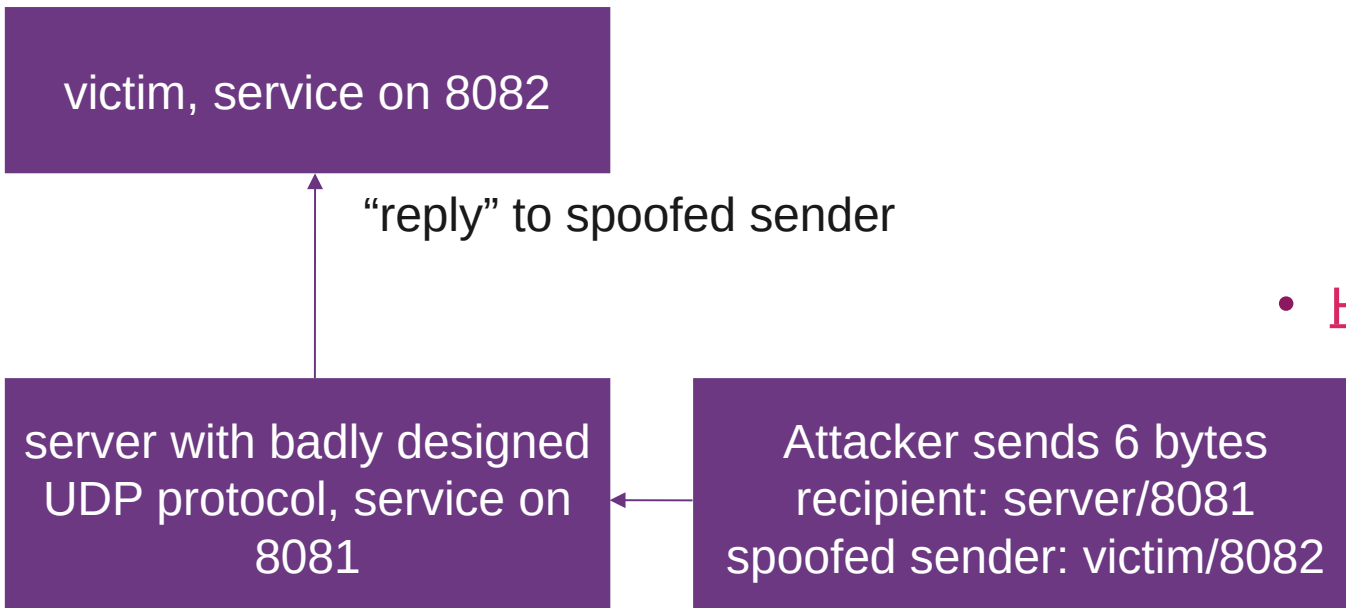
import (
    "net"
)

func main() {
    srv, _ :=
net.ResolveUDPAddr("udp", "127.0.0.1:8081")
    local, _ := net.ResolveUDPAddr("udp",
"127.0.0.1:0")
    conn, _ := net.DialUDP("udp", local, srv)
    defer conn.Close()
    conn.Write([]byte("5Anybody there?"))
}
```

- nc -u localhost 8081

Layer 4 - Transport

- [DDoS Amplification Attacks](#)
 - Request 10 bytes, reply 100 bytes → factor 10
- Local demo with server-ra/victim, and hping3
 - `hping3 --udp IP -p 8081 -E test.tmp -d 6 -s 8082 -c 1`



- Attacker in go/Java/node/c#
 - You need to spoof UDP packets, typically not supported in those languages
 - Go: `func DialUDP(network string, laddr, raddr *UDPAddr) (*UDPConn, error)`
 - laddr: we need to set here the victims IP/port
 - But go tries to bind to that
 - Not yours: “bind: cannot assign requested address”
- [Hping3](#): Pen test tool

hping3 is a command-line oriented IP, TCP, UDP, ICMP and RAW-IP packet assembler

Comparison – Transport Layer

TCP *

UDP *

SCTP *

(QUIC) *

- | | | | |
|--------------------------------|-------------------------------|-------------------------------|---|
| ■ Transport layer | ■ Transport layer | ■ Transport layer | ■ Transport layer* |
| ■ Connection oriented | ■ Connection less | ■ Connection oriented | ■ Connection oriented |
| ■ Reliable transfer | ■ Unreliable transfer | ■ Reliable transfer | ■ Reliable transfer |
| ■ Streams | ■ Messages | ■ Messages | ■ Multistream |
| ■ Guaranteed order | ■ Unordered | ■ User can choose | ■ Guaranteed order |
| ■ Widely used – HTTP/1, HTTP/2 | ■ Widely used – DNS, HTTP/3 | ■ WebRTC | ■ HTTP/3 |
| ■ Flow and congestion control | ■ No flow, congestion | ■ Flow and congestion control | ■ Flow and congestion |
| ■ Heavyweight | ■ Lightweight | ■ Heavyweight | ■ Heavyweight* |
| ■ Header size is 20 bytes | ■ Header size is 8 bytes | ■ Common header is 12 bytes | ■ Long: 17, short: 13*
(no sec) - 1200 |
| ■ Error checking and recovery | ■ Error checking, no recovery | ■ Error checking and recovery | ■ Integrity check |